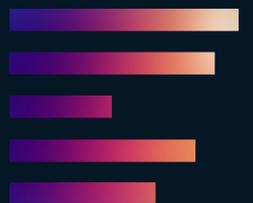


FALL 2024

MATH 230:
Abstract & Discrete
Mathematics

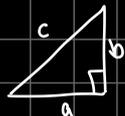


Pythagorean Theorem:

Friday, August 23

Generalized Pythagorean Theorem:

$$a^2 + b^2 = c^2$$

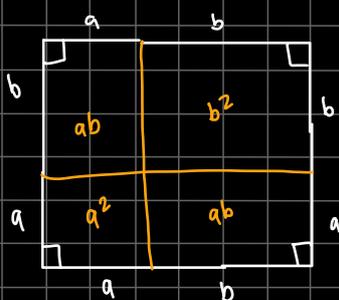
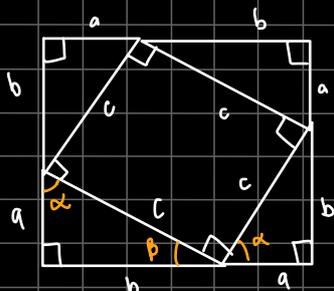


side length opposite the angle between a & b

Theorem: given a right triangle with side lengths a , b and c , where c is the length of the hypotenuse.

$$a^2 + b^2 = c^2$$

longest side



$$A = (a+b)^2 = a^2 + 2ab + b^2$$

$$A = 4\left(\frac{ab}{2}\right) + c^2 = 2ab + c^2$$

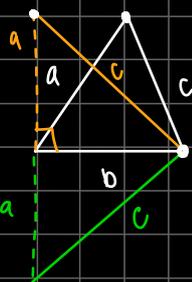
$$a^2 + 2ab + b^2 = 2ab + c^2$$

$$\therefore a^2 + b^2 = c^2$$

Converse of Pythagorean Theorem:

- Suppose we have a triangle with side lengths a , b , and c such that $a^2 + b^2 = c^2$.

- Then, the triangle is a right triangle.



$$a^2 + b^2 = c^2$$

$$c^2 = a^2 + b^2 = e^2$$

$$\therefore c = e$$

"If A then B":

- only value is A is true
- if A is false, we have a **vacuous truth**.

- Pythagoras' Theorem is both an

$$\text{"if A, then B"} \iff A \Rightarrow B$$

$$\text{"if B, then A"} \iff B \Rightarrow A$$

$$\therefore \text{"A if and only if B"} \iff A \iff B$$

- Suppose we have a triangle \triangle with side lengths a , b , and c , with c being the longest of the three.
- Then \triangle is a right triangle if and only if $a^2 + b^2 = c^2$.

Find a general solution to

$$\left. \begin{aligned} ax^2 + bx + c &= 0 \\ \text{when } a &\neq 0 \end{aligned} \right\} x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

Since $a \neq 0$, we can divide both sides by a .

$$\begin{aligned} \therefore \frac{ax^2 + bx + c}{a} &= \frac{0}{a} & \left(x + \frac{b}{2a}\right)^2 &= x^2 + \frac{b}{a}x + \frac{b^2}{4a^2} \\ \underbrace{x^2 + \frac{b}{a}x + \frac{c}{a}}_{\text{complete the square}} &= 0 & \therefore \left(x + \frac{b}{2a}\right)^2 + \frac{c}{a} &= \frac{b^2}{4a^2} \\ & & \left(x + \frac{b}{2a}\right)^2 &= \frac{b^2}{4a^2} - \frac{c}{a} \end{aligned}$$

$$\left(x + \frac{b}{2a}\right)^2 = \frac{b^2 - 4ac}{4a^2}$$

take the square root

$$x + \frac{b}{2a} = \pm \sqrt{\frac{b^2 - 4ac}{4a^2}} = \pm \frac{\sqrt{b^2 - 4ac}}{2a}$$

$$\therefore x = -\frac{b}{2a} \pm \frac{\sqrt{b^2 - 4ac}}{2a}$$

Law of Cosines:



$$\text{Then } c^2 = a^2 + b^2 - 2ab \cos \theta$$

$$\cos \theta = \frac{m}{a} \quad - m = a \cos \theta$$

$$- m = b - n \quad - b = n + m$$

$$- n = b - m$$

By pythagoras:

$$m^2 + n^2 = a^2$$

$$n^2 + h^2 = c^2$$

$$\therefore c^2 = n^2 + (a^2 - m^2)$$

$$= (n^2 - m^2) + a^2$$

$$= (n - m)(n + m) + a^2$$

$$= b(n - m) + a^2$$

$$= b(b - 2m) + a^2$$

$$= b^2 - 2bm + a^2$$

$$= a^2 + b^2 - 2ab \cos \theta$$

5. Proof

Monday, 26 August

Proof Template 1 Direct proof of an if-then theorem.

- Write the first sentence(s) of the proof by restating the hypothesis of the result. Invent suitable notation (e.g., assign letters to stand for variables).
- Write the last sentence(s) of the proof by restating the conclusion of the result.
- Unravel the definitions, working forward from the beginning of the proof and backward from the end of the proof.
- Figure out what you know and what you need. Try to forge a link between the two halves of your argument.

definition: An integer n is **even** if there exists an integer k such that $n = 2k$.

theorem: The sum of two even integers is even

proof: - let x, y be even.

- we want to show that $x+y$ is even

- since x, y are even, there exist m, n integers such that:

$$\begin{aligned} x &= 2m \\ y &= 2n \\ \left. \begin{array}{l} x = 2m \\ y = 2n \end{array} \right\} x+y &= 2c \quad \left. \begin{array}{l} \text{the sum of two} \\ \text{integers is} \end{array} \right\} \text{an integer} \\ x+y &= 2(m+n) \\ \text{let } c &= m+n \\ \therefore x+y &= 2c \quad \text{where } c \text{ is an integer} \\ \therefore x+y &\text{ is even } \quad \square \end{aligned}$$

definition: Let a, b be integers. We say $a \mid b$ ("a divides b") if there exists an integer k such that $b = ak$.

↓ allows $\frac{a}{b}$

theorem: If a, b, d are integers such that $d \mid a$ and $d \mid b$, then $d \mid a+b$.

proof: - Since $d \mid a$ and $d \mid b$, there exist integers m and n such that:

$$a = dm$$

$$b = dn$$

- Therefore $a+b = d(m+n)$

$$\text{let } c = m+n$$

$$\therefore a+b = dc$$

- c is an integer and $a+b = dc$

$$\therefore d \mid a+b \quad \square$$

definition: An integer n is **odd** if there exists an integer k such that $n = 2k+1$

theorem: - let n be an integer.

- n is even if and only if $n+1$ is odd.

proof: WTP: "If n is even, $n+1$ is odd" (\Rightarrow) sufficient
WTP: "If $n+1$ is odd, n is even" (\Leftarrow) necessary

(\Rightarrow): - Suppose n is even so $n = 2k$ for some integer k .

- Then $n+1 = 2k+1$, so $n+1$ is odd.

(\Leftarrow): - Suppose $n+1$ is odd so $n+1 = 2k+1$ for some integer k .

- Then $n = (2k+1)-1 = 2k$, so n is even. \square

definition: An integer $p > 1$ is **prime** if the only divisors of p are 1 and p

proposition: If $n > 1$ is an integer, then n^2+1 is not prime

$$\begin{aligned} \text{proof: } n^2+1 &= (n+1)(n^2-n+1) \\ \text{so } n+1 &\mid n^2+1 \end{aligned}$$

$$\text{NB: } n+1 \neq 1, n+1 \neq n^2+1$$

$$n > 1 \text{ so } n+1 > 1$$

$$n > 1 \text{ so } n^2 > n \text{ so } n^2+1 > n+1$$

$$\therefore n+1 \neq 1, n+1 \neq n^2+1 \quad \square$$

definition: A positive integer a is called **composite** provided there is an integer b such that $1 < b < a$ and $b \mid a$.

proposition: Let x be an integer. If $x > 1$, then x^2+1 is composite.

proof:

1. Let x be an integer and suppose that $x > 1$.

2. Note that $x^2+1 = (x+1)(x^2-x+1)$

3. Since x is an integer, both $(x+1)$ and (x^2-x+1) are integers.

4. Therefore $x+1 \mid x^2+1$

- Since $x > 1$, $x+1 > 1+1 = 2 > 1$.

- Also, since $x > 1$, $x < x^2$. Adding 1 to both sides, $x+1 < x^2+1$

- Thus, $x+1$ is an integer with $1 < x+1 < x^2+1$.

5. Since $1 < x+1 < x^2+1$ and $x+1 \mid x^2+1$, x^2+1 is composite. \square

6. Counterexample

Proof Template 3 Refuting a false if-then statement via a counterexample.

To disprove a statement of the form "If A, then B":
Find an instance where A is true but B is false.

disprove: If $a|b$ and $b|a$, then $a=b$.

- proof:
- Since $a|b$, $b=ka$ for some integer k .
 - Likewise, since $b|a$, $a=mb$ for some integer m .
 - Substituting b from step 1 into the equation in step 2: $a=m(ka)$
 $a=(mk)a$
 - If $a \neq 0$, $mk=1$
 $\therefore m=k=1$ or $m=k=-1$
 - If $m=k=1$, $a=(1)b$, $a=b$
 - If $m=k=-1$, $a=(-1)b$, $a=-b$
 - Therefore $a=\pm b$ □

disprove: If a and b are integers with $a|b$, $a \leq b$.

- proof: NTS: $a|b$ but $a > b$
- Let $a=2$, $b=0$
 - First, we show that $a|b$: $2|0$ is true because there exists an integer $k=0$ such that $b=ka$: $0=(0)2$
 - However $a \leq b$ does not hold: $2 \not\leq 0$ because $2 > 0$. □

disprove: If p and q are prime, then $p+q$ is composite.

- proof:
- Let $p=2$, $q=3$
 - First, we show that p and q are prime.
 - 2 is prime as its only factors are 1 and 2.
 - 3 is prime as its only factors are 1 and 3.
 - Now we evaluate $p+q$:
 - $p+q = 2+3=5$
 - We need to show that 5 is not composite (i.e., 5 is prime).
 - The factors of 5 are only 1 and 5
 - therefore, 5 is prime, not composite. □

7. Boolean Algebra

Wed, Aug 28

Proof Template 4 Truth table proof of logical equivalence

To show that two Boolean expressions are logically equivalent:
Construct a truth table showing the values of the two expressions for all possible values of the variables.
Check to see that the two Boolean expressions always have the same value.

theorem: $x \vee \neg x = T$

" \vee ": or
" \wedge ": and
" \neg ": not

proof:

x	$\neg x$	$x \vee \neg x$
T	F	T
F	T	T

□

OR

proof: - If x is true, then $\neg x$ is false, so $x \vee \neg x = T$.
- If x is false, then $\neg x$ is true, so $x \vee \neg x = T$. □

" \rightarrow ": implies

" \leftrightarrow ": if & only if

x	y	$x \rightarrow y$	x	y	$x \leftrightarrow y$
T	T	T	T	T	T
T	F	F	T	F	F
F	T	T	F	T	F
F	F	T	F	F	T

} vacuously true

proposition: $(x \rightarrow y) \wedge (y \rightarrow x) = x \leftrightarrow y$

proof:

x	y	$x \rightarrow y$	$y \rightarrow x$	$(x \rightarrow y) \wedge (y \rightarrow x)$
T	T	T	T	T
T	F	F	T	F
F	T	T	F	F
F	F	T	T	T

□

Theorem 7.2:

$$\begin{aligned} \textcircled{1} \quad & x \wedge y = y \wedge x \\ & x \vee y = y \vee x \end{aligned} \quad \left. \vphantom{\begin{aligned} \textcircled{1} \quad & x \wedge y = y \wedge x \\ & x \vee y = y \vee x \end{aligned}} \right\} \text{commutative properties}$$

$$\begin{aligned} \textcircled{2} \quad & (x \wedge y) \wedge z = x \wedge (y \wedge z) \\ & (x \vee y) \vee z = x \vee (y \vee z) \end{aligned} \quad \left. \vphantom{\begin{aligned} \textcircled{2} \quad & (x \wedge y) \wedge z = x \wedge (y \wedge z) \\ & (x \vee y) \vee z = x \vee (y \vee z) \end{aligned}} \right\} \text{associative properties}$$

$$\begin{aligned} \textcircled{3} \quad & \overbrace{x \wedge T}^{x \times 1} = x \\ & \underbrace{x \vee F}_{x + 0} = x \end{aligned} \quad \left. \vphantom{\begin{aligned} \textcircled{3} \quad & \overbrace{x \wedge T}^{x \times 1} = x \\ & \underbrace{x \vee F}_{x + 0} = x \end{aligned}} \right\} \text{identity elements}$$

$$\textcircled{4} \quad \neg(\neg x) = x \quad \left. \vphantom{\textcircled{4}} \right\} \text{involution}$$

$$\begin{aligned} \textcircled{5} \quad & x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z) \\ & x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z) \end{aligned} \quad \left. \vphantom{\begin{aligned} \textcircled{5} \quad & x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z) \\ & x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z) \end{aligned}} \right\} \text{distributive properties}$$

$$\textcircled{6} \quad x \wedge (\neg x) = F, \quad x \vee (\neg x) = T$$

$$\begin{aligned} \textcircled{7} \quad & \neg(x \wedge y) = (\neg x) \vee (\neg y) \\ & \neg(x \vee y) = (\neg x) \wedge (\neg y) \end{aligned} \quad \left. \vphantom{\begin{aligned} \textcircled{7} \quad & \neg(x \wedge y) = (\neg x) \vee (\neg y) \\ & \neg(x \vee y) = (\neg x) \wedge (\neg y) \end{aligned}} \right\} \text{DeMorgan's Laws}$$

proposition: $x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z)$

proof:

x	y	z	$y \wedge z$	$x \vee (y \wedge z)$	$x \vee y$	$x \vee z$	$(x \vee y) \wedge (x \vee z)$
T	T	T	T	T	T	T	T
T	T	F	F	T	T	T	T
T	F	T	F	T	T	T	T
T	F	F	F	T	T	T	T
F	T	T	T	T	T	T	T
F	T	F	F	F	T	F	F
F	F	T	F	F	F	T	F
F	F	F	F	F	F	F	F

OR

proof: - if $x = T$, then $x \vee (y \wedge z) = T$,
likewise both $(x \vee y)$ and $(x \vee z) = T$.
- so, $(x \vee y) \wedge (x \vee z) = T$
- if $x = F$, then $x \vee (y \wedge z) = y \wedge z$
likewise $x \vee y = y$
 $x \vee z = z$
- so, $(x \vee y) \wedge (x \vee z) = y \wedge z$ \square

8. Lists

Fri, Aug 30

definition: A list is an ordered collection of elements in a sequence.

$(1, 2, 3), ((1, 2), (1, 2, 3))$

$(1, 2, 3) \neq (3, 2, 1)$

1. How many lists with 2 elements exist if the elements are $(1, 2, 3, 4)$.

$4 \times 4 = 16$ //

Multiplication Principle: The number of list with 2 elements where the first element is one of m things and the second element is one of n things is $m \times n$.

Ex: count the divisors of 10^n
if $d \mid 10^n$, then $d = 2^a \cdot 5^b$
where $0 \leq a \leq n$
and $0 \leq b \leq n$

The # of such divisors is the number of lists of the form (a, b) where $0 \leq a \leq n$
 $0 \leq b \leq n$

answer: $(n+1)(n+1) = (n+1)^2$

- A list with k elements and each element is one of n possibilities.
- # of lists = n^k .
- What if we don't allow repeats?

$n(n-1)(n-2) \dots (n-k+1) = (n)_k$

Ex: 100m final, 8 lanes, 3 winners

of possible podiums = $8(8-1)(8-2)$
 $= 8 \times 7 \times 6$

$n! = n(n-1)(n-2) \dots (1)$ for $n \geq 1$

\downarrow
 $= \frac{n!}{(n-k)!}$ if $0 \leq k \leq n$

10. Sets I: Introduction, Subsets

definition: A set is an unordered collection of elements.

$\{1, 2\}, \{1, 2\}, \{1, 2, 3\}, \{1, 2, 3\}$

$\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$

$\mathbb{N} = \{1, 2, 3, \dots\} = \mathbb{Z}^+$

$\mathbb{N}_0 = \{0, 1, 2, 3, \dots\}$

$\{1, 2, 3\} = \{3, 2, 1\}$

- Two sets are equal if they have the same elements.

$A = \{(a, b, c) \mid a, b, c \in \mathbb{Z} \text{ and } a^2 + b^2 = c^2\}$

$B = \{(a, b, c) \mid \begin{matrix} a = x^2 - y^2, & b = 2xy, \\ c = x^2 + y^2 \end{matrix} \text{ for some } x, y \in \mathbb{Z}\}$

$C = \{\frac{a}{b} \mid a, b \in \mathbb{Z} \text{ and } b \neq 0\}$

$D = \{\frac{a}{b} \mid a \in \mathbb{Z}, b \in \mathbb{N}\}$

Suppose: $(a, b, c) \in B$,
then, there exist $x, y \in \mathbb{Z}$ s.t.

$a = x^2 - y^2$

$b = 2xy$

$c = x^2 + y^2$

Note: $a, b, c \in \mathbb{Z}$

$\therefore a^2 + b^2 = (x^2 - y^2)^2 + (2xy)^2$
 $= x^4 - 2x^2y^2 + y^4 + 4x^2y^2$
 $= x^4 + 2x^2y^2 + y^4 = (x^2 + y^2)^2 = c^2$

- since $a, b, c \in \mathbb{R}$ and $a^2 + b^2 = c^2$

$$(a, b, c) \in A$$

$$\therefore B \subseteq A$$

- However, $A \not\subseteq B$.

- For example, $(4, 3, 5) \in A$

- For $(4, 3, 5) \in B$, we need $3 = x^2 + y^2$ for some $x, y \in \mathbb{R}$.

- However $x^2 + y^2 = \frac{2}{3} \notin \mathbb{R}$.

- given 2 sets S and T .

To prove $S = T$.

- You need:

$$(1) S \subseteq T \quad (s \in S, s \in T)$$

$$(2) T \subseteq S \quad (t \in T, t \in S)$$

- show that $C = D$

- suppose $d \in D$:

- so $d = \frac{a}{b}$ with $a \in \mathbb{Z}$ and $b \in \mathbb{N}$

- so $a, b \in \mathbb{Z}$ and $b \neq 0$

- hence $\frac{a}{b} \in \mathbb{C}$

- so $D \subseteq C$.

- suppose $c \in C$:

- so $c = \frac{a}{b}$ with $a, b \in \mathbb{Z}$ and $b \neq 0$

- so $a \in \mathbb{Z}, b \in \mathbb{Z} \wedge b \neq 0$

- if $b \in \mathbb{N}$, then $\frac{a}{b} \in D$

- if $b \notin \mathbb{N}$, b is a negative integer bc $b \neq 0$

$$\frac{a}{b} = \frac{a(-1)}{b(-1)} = \frac{-a}{-b}$$

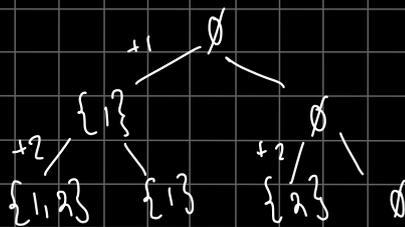
- $b \in \mathbb{N}$ and $b \neq 0$ so $-b \in \mathbb{N}$

- $-a \in \mathbb{Z}$

- $\therefore \frac{-a}{-b} \in D$, so $C \subseteq D$.

$$A = \{1, 2\}$$

$$\text{Subsets: } \emptyset, \{1\}, \{2\}, \{1, 2\}$$



$$A = \{1, 2, 3\}$$

$$\text{Subsets: } \emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}$$

$$A = \{1, 2, 3, \dots, n\}$$

$$\text{Subsets} = 2^n$$

↓

definition: the power set of A is the sets of all sets of A

$$\text{pow}(A) = 2^A$$

11. Quantifiers

Wed, Sep 4

- \exists - there exists - $!$ - unique

- Existential statement: to prove them, you need an example

- Example: There is an even integer that is prime.

- Proof: 2 is even and 2 is prime

- \forall - for all \Rightarrow Proof: Negate the prove

- Example: The sum of two even integers is an even integer.

- Example: $A \subseteq B \Rightarrow \forall x \in A, \forall x \in B$

"Every Integer is prime"

logic: $\forall n \in \mathbb{Z}, n$ is prime.

negate: $\exists n \in \mathbb{Z}, n$ is not prime

sentence: "There is an integer that is not prime"

"There is an integer x such that for all integers y , $x+y=0$ "

- $\exists x \in \mathbb{Z}, \forall y \in \mathbb{Z}, x+y=0$
- $\forall x \in \mathbb{Z}, \exists y \in \mathbb{Z}, x+y \neq 0$
- "For all integers x , there exists an integer y such that $x+y \neq 0$."

"Some integer is divisible by 7"

- $\exists n \in \mathbb{Z}, 7|n$
- $\forall n \in \mathbb{Z}, \exists x \in \mathbb{Z}, 7 \nmid xn$
- "All integers are not divisible by 7"

"There is no line that goes through all 3 line segments (unless it goes through a vertex)"

Negation: "There is a line that goes through all 3 line segments without going through a vertex."

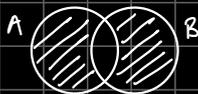
12. Sets II: Operations

union: $A \cup B = \{x \mid x \in A \text{ or } x \in B\}$

intersection: $A \cap B = \{x \mid x \in A \text{ and } x \in B\}$

difference: $A \setminus B = A - B = \{x \mid x \in A \text{ and } x \notin B\}$

sym. diff.: $A \Delta B = (A - B) \cup (B - A)$



Cartesian product: $A \times B = \{(x,y) \mid x \in A \text{ and } y \in B\}$

De Morgan's Laws
 $A - (B \cup C) = (A - B) \cap (A - C)$
 $A - (B \cap C) = (A - B) \cup (A - C)$

Theorem: If A and B are finite and $|A|=m$ and $|B|=n$, then $|A \times B| = m \cdot n$

$$\text{NTS: } A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$$

1. Let $x \in A \cup (B \cap C)$

- so $(x \in A) \text{ or } (x \in B \text{ and } x \in C)$
- if $(x \in A) \text{ or } (x \in B)$, $x \in (A \cup B)$
 - if $(x \in A) \text{ or } (x \in C)$, $x \in (A \cup C)$
- $\therefore x \in (A \cup B) \cap (A \cup C)$

2. Let $y \in (A \cup B) \cap (A \cup C)$

- so $y \in (A \cup B)$ and $y \in (A \cup C)$
- if $y \in A$, $y \in A$ for both conditions
 - if $y \notin A$, $y \in B$ and $y \in C$
 - if $y \in B$ and $y \in C$, $y \in B \cap C$
- $\therefore y \in A \cup (B \cap C)$

Theorem (Principle of Inclusion-Exclusion):

Suppose A, B are finite sets:

$$|A \cup B| = |A| + |B| - |A \cap B|$$

Proof: - Let a_1, a_2, \dots, a_n be the elements of A not in B .
 - Let b_1, b_2, \dots, b_m be the elements of B not in A .
 - Let c_1, c_2, \dots, c_r be the elements of $A \cap B$.

$$\begin{aligned} \text{Then } |A \cup B| &= n + m + r \\ |A| &= n + r \\ |B| &= m + r \end{aligned} \quad \therefore |A| + |B| - |A \cap B| = (n+r) + (m+r) - (m+r)$$

$$\begin{aligned} |A \cup B \cup C| &= |A \cup B| + |C| - |(A \cup B) \cap C| \\ &= |A| + |B| - |A \cap B| + |C| - |(A \cap C) \cup (B \cap C)| \\ &= |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C| \end{aligned}$$

How many positive integers ≤ 1000 are even, mult. of 3 or mult. of 5?

A: Even integers between 1 & 1000 (incl.)

B: Mult. of 3 between 1 & 1000

C: Mult. of 5 between 1 & 1000 (incl.)

$$\begin{aligned} |A \cup B \cup C| &= |A| + |B| + |C| - \underbrace{|A \cap B|}_{\text{mult. of 6}} - \underbrace{|A \cap C|}_{\text{mult. of 10}} - \underbrace{|B \cap C|}_{\text{mult. of 15}} + |A \cap B \cap C| \\ &= 500 + 333 + 200 - 166 - 100 - 66 + 33 \end{aligned}$$

Combinations with repetition
 Dozen donuts, 3 flavors, how many ways?

$$n=2$$

$$r=3$$



14 choices for the 2 separators $\binom{14}{2}$

14. Relations

Mon, Sep 9

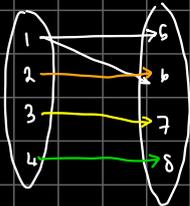
A relation R on $A \times B$ is a subset of $A \times B$.

Ex: $A = \{1, 2, 3, 4\}$
 $B = \{5, 6, 7, 8\}$

$$R = \{(1,5), (1,6), (2,6), (3,7), (4,7)\}$$

"1 is related to 5"
 "1 R 5"
 " $(1,5) \in R$ "
 if $(x,y) \in R, xRy$.

- 1R5, 1R6, 2R6, 3R7, 4R7
 1R7, 1R8, 2R5, 2R7, 2R8



SOME types of relations:

- Reflexive
- Irreflexive
- Transitive
- Symmetric
- Antisymmetric

Let R be a relation on A .

① Reflexive: $(=)$

- R is reflexive if $\forall a \in A, (a,a) \in R$ or aRa

② Irreflexive: No loops

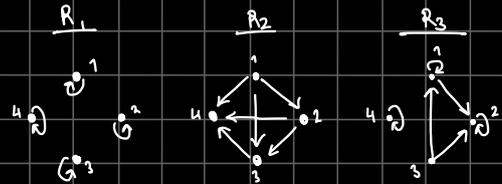
- R is irreflexive if $\forall a \in A, (a,a) \notin R$ or $a \not R a$

③ Symmetric: (\neq)

- R is symmetric if whenever $(a,b) \in R, (b,a) \in R$. $aRb \Rightarrow bRa$

④ Antisymmetric: (\leq) [any arrow goes 1 way]

- R is antisymmetric if $(a,b) \in R$ and $(b,a) \in R$ implies $a = b$
 $(aRb \wedge bRa) \Rightarrow a = b$



⑤ Transitive: $(<)$

- R is transitive if $(a,b) \in R$ and $(b,c) \in R$ implies $(a,c) \in R$. $(xRy \wedge yRz) \Rightarrow xRz$

R is a relation on a set A if $R \subseteq A \times A$

$$A = \{1, 2, 3, 4\}$$

1. Let R_1 be the relation "is equal to"

$$R_1 = \{(1,1), (2,2), (3,3), (4,4)\}$$

2. Let R_2 be the relation "is less than"

$$R_2 = \{(1,2), (1,3), (1,4), (2,3), (2,4), (3,4)\}$$

3. Let R_3 be the relation such that:

$$R_3 = \{(1,1), (1,2), (2,2), (3,1), (3,2), (4,1)\}$$

Theorem:

- Let R be a relation on a non-empty set A
- Then R is symmetric and antisymmetric if and only if $R \subseteq \{(a,a) \mid a \in A\}$

Proof:

(\Leftarrow) :

- Suppose $(a,b) \in R$
- Since $R \subseteq \{(a,a) \mid a \in A\}$
 $\therefore a = b$
 so $(b,a) \in R$
- Therefore R is symmetric.
- Suppose $(a,b) \in R$ and $(b,a) \in R$
- since $(a,b) \in R, a = b$
- So, it's antisymmetric.

⇒:

- Suppose R is symmetric and antisymmetric.
- Let $(a, b) \in R$.
- Since R is symmetric, $(b, a) \in R$.
- Since R is antisymmetric and $(a, b) \in R$ and $(b, a) \in R$,
so, $a = b$.
- Therefore $(a, a) \in \{ (x, x) \mid x \in A \}$

$$A = \{1, 2, 3, 4\}$$

- How many relations of A are symmetric and antisymmetric.

$$2^4 = 32 \text{ // (Power sets)}$$

- # of relations: 2^{n^2} (diagonals and off-diagonals)
- # of symmetric: $2^n \cdot 2^{\frac{n^2-n}{2}} = 2^{\frac{n^2+n}{2}}$
- # of antisymmetric: $2^n \cdot 3^{\frac{n^2-n}{2}}$
- # of reflexive: 2^{n^2-n} (3 choices when $a \neq b$)
- # of irreflexive: 2^{n^2-n} (include (a,b) but not (b,a) , include (a,a) but not (a,b) , exclude both)
- # of transitive: open problem.

15. Equivalence Relations

Wed, Sep 18

- We say a relation R on set A is an equivalence relation if R is reflexive, symmetric, and transitive.

$\{(1,1), (2,2), (3,3)\}$ on $\{1,2,3\}$ is an equiv. relation.

- Let R be the relation on \mathbb{Z} where a is related to b if they both have the same parity.

Prove it's an equivalence relation:

Reflexive: A number a has the same parity as itself—so $(a, a) \in R \therefore R$ is reflexive.

Symmetric: Suppose $(a, b) \in R$.
 case 1: a and b are both even, then $(b, a) \in R$.
 case 2: a and b are both odd, then $(b, a) \in R$.
 $\therefore R$ is symmetric.

Transitive: Suppose $(a, b) \in R$ and $(b, c) \in R$.
 Case 1: a is even
 Since $(a, b) \in R$, b is even
 Since $(b, c) \in R$, c is even.
 a and c have the same parity $\therefore (a, c) \in R$

Case 2: a is odd.

Since $(a, b) \in R$, b is odd.

Since $(b, c) \in R$, c is odd.

Since a and c have the same parity, $(a, c) \in R$.

$\therefore R$ is transitive.

→ equivalence relation

Congruence mod n :

they differ by a multiple of n .

- We say $a \equiv b \pmod{n}$ iff $n \mid (b-a)$
- a is congruent to $b \pmod{n}$, $a \equiv b \pmod{n} \Rightarrow a \pmod{n} = b \pmod{n}$

Consider the relation $R_n: \{ (a, b) \mid a \equiv b \pmod{n} \}$
 $\therefore (a, b) \in R_n$ if $n \mid b-a$

Prove equivalence relation on \mathbb{Z}

Reflexive: let $a \in \mathbb{Z}$

- NTS: $(a, a) \in R_n$, $a \equiv a \pmod{n}$

$$n \mid a-a$$

$$- a-a = 0, n \mid 0 \text{ because } 0 = 0 \cdot n$$

$$\therefore n \mid a-a$$

$$\therefore a \equiv a$$

- Thus, R_n is reflexive.

Symmetric: let $(a, b) \in R_n$

- NTS: $(b, a) \in R_n$, $b \equiv a \pmod{n}$, $n \mid a-b$, $a-b = kn$

- Since $(a, b) \in R_n$, $a \equiv b \pmod{n}$,

$$\text{so } n = b-a$$

- there $\exists k \in \mathbb{Z}$ s.t. $b-a = k \cdot n$

$$a-b = (-k)n$$

- Since $-k \in \mathbb{Z}$, $n \mid a-b \therefore (b, a) \in R_n$.

- Thus, R_n is symmetric.

Transitive: suppose $(a, b) \in R_n$ and $(b, c) \in R_n$

- NTS: $(a, c) \in R_n$, $a \equiv c \pmod{n}$, $n \mid c-a$.

- Since $(a, b) \in R_n$, $n \mid b-a$

- Since $(b, c) \in R_n$, $n \mid c-b$

- Since $n \mid b-a$ and $n \mid c-b$,
 $n \mid (b-a) + (c-b) = c-a$

- So, $(a, c) \in R_n$.. R_n is transitive.

$[2]_5 = \{ \dots, x-2m, x-m, x, x+m, x+2m, \dots \}$ equivalent class of $x \pmod{5} = [x]_5$

Ex: What is the equivalence class of 2 with respect to congruence modulo 5?

- $[2] = \{ x \in \mathbb{Z} : (x, 2) \in R \}$

$$\therefore x \equiv 2 \pmod{5} \Rightarrow x \pmod{5} = 2 \pmod{5} = 2$$

$$\therefore [2] = \{ \dots, -8, -3, 2, 7, 12, \dots \}$$

Equivalence Classes

- Let R be an equivalence relation on A .
- Let $a \in R$
- $[a] = \{x \in A : (x, a) \in R\}$
the subset of some equivalence relation R which includes all elements that are equivalent to each other.
- In the relation with parity, there are 2 equivalence classes.
 $[0]$ - set of even numbers
 $[1]$ - set of odd numbers

- Ex: $A = \{1, 2, 3, 4, 5\}$, $R = \{(a, b) : a+b \text{ is even}\}$
 $[1] = \{1, 3, 5\}$ $[3] = \{1, 3, 5\}$ $[5] = \{1, 3, 5\}$
 $[2] = \{2, 4\}$ $[4] = \{2, 4\}$
 $[1] = [3] = [5]$ $[2] = [4]$ \therefore there are only 2 equivalence classes

Theorem: Let R be an equivalence relation on A

① If $(a, b) \in R$, $[a] = [b]$

② If $[a] = [b]$, $(a, b) \in R$

③ $(a, b) \in R$, $[a] \cap [b] = \emptyset$

④ $\bigcup_{a \in A} [a] = A$

\downarrow union of all equivalence classes

Proof:

① $(a, b) \in R$ NTS: $[a] = [b]$

- (\Rightarrow) - Let $x \in [a]$ NTS: $[a] \subseteq [b]$
 - Then $(a, x) \in R$ $\rightarrow x \in [b]$
 - R is symmetric so $(x, a) \in R$ $\rightarrow (b, x) \in R$
 - $(x, a) \in R$ and $(a, b) \in R$
 so by transitivity, $(x, b) \in R$,
 by symmetry, $(b, x) \in R$.
 - So, $x \in [b]$
 - Thus $[a] \subseteq [b]$ NTS: $[b] \subseteq [a]$

- (\Leftarrow) - Let $y \in [b]$, so $(b, y) \in R$ $\rightarrow (a, y) \in R$
 - since $(a, b) \in R$ and $(b, y) \in R$, $(a, y) \in R$,
 - Hence, $y \in [a]$
 - So $[b] \subseteq [a]$
 therefore $[a] = [b]$.

② $[a] = [b]$ NTS: $(a, b) \in R \rightarrow b \in [a]$

- By reflexivity $a \in [a]$, $b \in [b]$.
- $b \in [b] \Rightarrow b \in [a]$ because $[a] = [b]$
- Since $b \in [a]$, $(a, b) \in R$ \square

② $\bigcup_{a \in A} [a] = A$

- Let $x \in A$, then $x \in [x]$
- so, $x \in \bigcup_{a \in A} [a]$
- therefore $A \subseteq \bigcup_{a \in A} [a]$
- $[a] = \{b \in A \mid (a, b) \in R\} \subseteq A$,
 so $\bigcup_{a \in A} [a] = A$

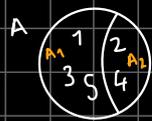
③ $(a, b) \notin R$, $[a] \cap [b] = \emptyset$

We'll prove the contrapositive ($x \rightarrow y = \neg y \rightarrow \neg x$)

- If $[a] \cap [b] \neq \emptyset$, $(a, b) \in R$.

- Let $x \in [a] \cap [b]$.
- so $(a, x) \in R$ and $(b, x) \in R$
- By symmetry, $(x, b) \in R$
- By transitivity, since $(a, x) \in R$ and $(x, b) \in R$,
 $(a, b) \in R$ \square

16. Partitions



- Let R be an equivalence relation on a set A .
- The equivalence classes of R form a partition of the set of A

(Partition) Let A be a set. A partition of (or on) A is a set of nonempty pairwise disjoint sets whose union is A .

- ① A **partition** is a set of sets: each member of a partition is a subset of A . The members are called **parts**.
- ② The parts of a partition are **nonempty**.
- ③ The parts of a partition are **pairwise disjoint**.
 No two sets of a partition may have any element in common.
 - if $[a] \neq [b]$, $[a] \cap [b] = \emptyset$
- ④ The union of the parts is the original set.
 $\bigcup [a] = A$
 $[a] \in A$

(Counting Equivalence Classes) Let R be an equivalence class on a finite set A . If all the equivalence classes of R have the same size, m , then the number of equivalence classes is $\frac{|A|}{m}$.

17. Binomial Coefficients

Fri, Sep 20

Pascal's Triangle	1						$n=0$
	1	1					$n=1$
	1	2	1				$n=2$
	1	3	3	1			$n=3$
	1	4	6	4	1		$n=4$
	1	5	10	10	5	1	$n=5$

(Binomial Coefficient): $\binom{n}{k} = nC_k$
 = "n choose k"

- number of ways of choosing k elements from a set of n elements (if the order doesn't matter)
- number of k-element subsets of an n-element set

(Pascal's Identity): $\binom{n}{k} + \binom{n}{k+1} = \binom{n+1}{k+1}$

$\frac{n(n-1)(n-2)\dots(n-k+1)}{k(k-1)(1)\dots(1)} \rightarrow$ enforces an order
 \rightarrow to remove ordering

$$\therefore \binom{n}{k} = \frac{n!}{k!(n-k)!} = \frac{n(n)}{k!}$$

Why is it true that:

$$(a+b)^n = \sum_{k=0}^n \binom{n}{k} a^{n-k} b^k \quad \text{Newton's Binomial Theorem}$$

Ex: $(a+b)^5 = (a+b)(a+b)(a+b)(a+b)(a+b)$
 \downarrow
 two choices when expanding
 $\therefore 2^5 = 32$
 \downarrow
 sum of row $n=5$ of Δ

Proving Pascal's Identity:

- Let n, k be nonnegative integers with $k+1 \leq n$

$$\binom{n}{k} + \binom{n}{k+1} = \binom{n+1}{k+1}$$

Algebraic Proof:

$$\begin{aligned} \binom{n}{k} + \binom{n}{k+1} &= \frac{n!}{k!(n-k)!} + \frac{n!}{(k+1)!(n-k-1)!} \\ &= \frac{n!(k+1)! + n!(n-k)!}{(k+1)!(n-k)!} \\ &= \frac{n!(n+1)}{(k+1)!(n+1-(k+1))!} \\ &= \frac{(n+1)!}{(k+1)!(n+1-(k+1))!} = \binom{n+1}{k+1} \quad \square \end{aligned}$$

Combinatorial Proof:

- Let's count the number of subsets of $\{1, 2, \dots, n+1\}$ with $k+1$ elements
 - ① it's $\binom{n+1}{k+1}$
 - ② (consider whether the subsets have $n+1$ or not.
 - there are $\binom{n}{k}$ subsets that contain $n+1$.
 - there are $\binom{n}{k+1}$ subsets that don't contain $n+1$.

(Theorem) Let $0 \leq k \leq n$ are integers
 then $\binom{n}{k} = \binom{n}{n-k}$

Proof: Algebraic

$$\binom{n}{k} = \frac{n!}{k!(n-k)!} \quad \binom{n}{n-k} = \frac{n!}{(n-k)!k!}$$

same! \leftarrow

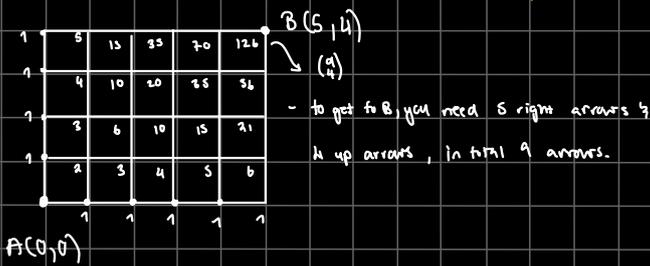
Proof: Combinatorial

- If you choose a subset with k elements of $\{1, 2, \dots, n\}$
- It's complement is a subset with $n-k$ elements.

e.g: the # of subsets of $\{1, 2, 3, 4, 5\}$ with 2 elements = # of subsets with 3 elements.

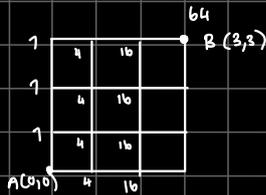
Counting Lattice Paths:

① How many paths from $A \rightarrow B$ if you can only move $\rightarrow \uparrow$? N-E lattice path



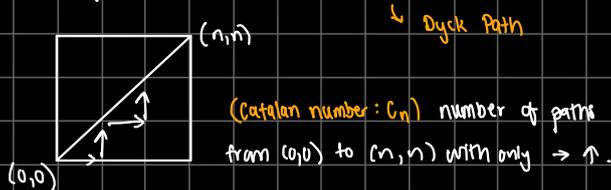
- In general, the number of paths from $(0,0)$ to (n,k) is $\binom{n+k}{n}$ or $\binom{n+k}{k}$ \rightarrow this is because once we choose the N's, the E's are given, and vice versa.

② Now what if we allow $\rightarrow \uparrow \downarrow$, but you can't repeat a vertex?



- In general: $(b+1)^a$

③ Now, you can't cross the diagonal, only $\rightarrow \uparrow$.



* Note, you can never have more \uparrow 's than \rightarrow 's.

$$\frac{1}{n+1} \binom{2n}{n}$$

Poker Hands:

- Deck has 52 numbers, 13 numbers, 4 suits.

- Pair: same number

- 3-of-a-kind: same number

- 4-of-a-kind: same number

- Full House: 3-of-a-kind + Pair

- Straight: 5 consecutive numbers

- Flush: 5 in same suit

- Straight Flush: Straight + Flush

1. How many poker hands?

$$\binom{52}{5}$$

2. How many poker hands are there with just one pair?

$$\binom{13}{1} \binom{4}{2} \binom{12}{3} \binom{4}{1} \binom{4}{1} \binom{4}{1}$$

pair last three

3. How many poker hands are there with two pairs?

$$\binom{13}{2} \binom{4}{2} \binom{4}{2} \binom{11}{1} \binom{4}{1}$$

two pair last one

4. How many poker hands are there with three of a kind? \uparrow and not a full house

$$\binom{13}{1} \binom{4}{3} \binom{12}{2} \binom{4}{1} \binom{4}{1}$$

three of a kind last two

5. How many poker hands are there with four of a kind?

$$\binom{13}{1} \binom{4}{4} \binom{12}{1} \binom{4}{1}$$

four of a kind last one

6. How many poker hands are a full house?

$$\binom{13}{1} \binom{4}{3} \binom{12}{1} \binom{4}{2}$$

three of a kind pair

7. How many poker hands are a straight flush?

$$10 \times \binom{4}{1} = 40 //$$

\downarrow from A 2 3 4 5

\vdots

10 J Q K A

8. How many poker hands are a flush? \uparrow but not a straight

$$\binom{13}{5} \binom{4}{1} = 40$$

9. How many poker hands are a straight? \uparrow but not a flush

$$10 \times \binom{4}{1}^5 = 40$$

\uparrow suit selection

10. How many poker hands are a royal flush?

$$\binom{4}{1} \binom{5}{5} = 4 //$$

suit royal flush

2.2. Induction

- We have some statement that depends on a variable n which is a positive integer.
 $S(n)$.

- NTS: $S(n)$ is true $\forall n \geq n_0$.

- Steps:

1. Prove $S(n_0)$ is true. → Base case

2. Prove $S(k)$ implies $S(k+1)$ → Inductive step

- You can conclude that it's true $\forall n \geq n_0$.

Ex: Prove $1+2+\dots+n = \frac{n(n+1)}{2}$ for some $n \geq 1$

Base case: $n=1$

$$1 = \frac{1(1+1)}{2} = \frac{2}{2} = 1 \quad \square$$

Inductive hypothesis:

- Suppose $1+2+\dots+k = \frac{k(k+1)}{2}$ for some $k \geq 1$. (1)

Inductive step: (NTS) $1+2+\dots+k+(k+1) = \frac{(k+1)(k+2)}{2}$

- Adding $k+1$ to both sides of (1):

$$\begin{aligned} 1+2+\dots+k+(k+1) &= \frac{k(k+1)}{2} + (k+1) \\ &= \frac{k^2+k+2k+2}{2} \\ &= \frac{k^2+3k+2}{2} = \frac{(k+1)(k+2)}{2} \quad \square \end{aligned}$$

Ex: Hockey sticks identity

$$\binom{r}{r} + \binom{r+1}{r} + \dots + \binom{n}{r} = \binom{n+1}{r+1} \quad \text{for } n \geq r.$$

Base case: $n=r$

$$\binom{r}{r} = \binom{r+1}{r+1} \Rightarrow 1 = 1$$

Induction hypothesis:

- Suppose $\binom{r}{r} + \binom{r+1}{r} + \dots + \binom{k}{r} = \binom{k+1}{r+1}$

Inductive case:

$$\text{NTS: } \binom{r}{r} + \binom{r+1}{r} + \dots + \binom{k}{r} + \binom{k+1}{r} = \binom{k+2}{r+1}$$

$$\therefore \binom{k+1}{r+1} + \binom{k+1}{r} = \binom{k+2}{r+1} \quad \square$$

Theorem: Let $n \in \mathbb{Z} : n \geq 0$

Then $2^n \geq 1+n$

Base case: $n=0$

$$2^0 \geq 1+0$$

$$1 \geq 1$$

Inductive hypothesis: Suppose $2^k \geq 1+k$

Inductive case: $2^{k+1} \geq k+2$

$$\therefore 2^k \cdot 2 \geq k+2$$

$$(k+1)2 \geq k+2$$

$$2k+2 \geq k+2 \quad \forall k \geq 0 \quad \square$$

Prove: $\frac{1}{1} + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{2^n} \geq 1 + \frac{n}{2}$, for $n \geq 0$

Base case: $n=0$

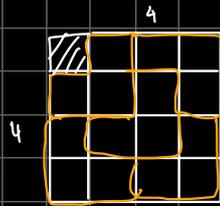
$$\frac{1}{2^0} \geq 1 + \frac{0}{2} \Rightarrow 1 \geq 1$$

Inductive hypothesis: Suppose $\frac{1}{1} + \frac{1}{2} + \dots + \frac{1}{2^k} \geq 1 + \frac{k}{2}$

Inductive case: $\frac{1}{1} + \frac{1}{2} + \dots + \frac{1}{2^k} + \frac{1}{2^{k+1}} \geq 1 + \frac{k+1}{2} \geq 1 + \frac{k}{2} + \frac{1}{2}$

$$1 + \frac{k}{2} + \frac{1}{2^{k+1}} \geq 1 + \frac{k}{2} + \frac{1}{2}$$

$$\frac{1}{2^{k+1}} \geq \frac{1}{2} \quad \forall k \geq 0 \quad \square$$



Theorem: Let $n \geq 1$. A $2^n \times 2^n$ grid with a unit square removed can be tiled with L-trominoes.

Base case: $n=1$



Inductive hypothesis: Suppose any $2^k \times 2^k$ grid with a square removed can be tiled with L-trominoes.

Inductive step:



Prove:

$$f_1 = 1$$

$$f_2 = 1$$

$$f_n = f_{n-1} + f_{n-2} \quad \text{for } n \geq 3$$

1, 1, 2, 3, ...

$$f_n = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^n - \left(\frac{1-\sqrt{5}}{2}\right)^n}{\sqrt{5}}$$

Base case: $n=1$

$$f_1 = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^1 - \left(\frac{1-\sqrt{5}}{2}\right)^1}{\sqrt{5}} = \frac{\sqrt{5}}{\sqrt{5}} = 1$$

$n=2$

$$f_2 = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^2 - \left(\frac{1-\sqrt{5}}{2}\right)^2}{\sqrt{5}} = \frac{6+2\sqrt{5}}{4} - \frac{6-2\sqrt{5}}{4}{\sqrt{5}} = \frac{4\sqrt{5}}{4\sqrt{5}} = 1$$

Inductive hypothesis:

Suppose $f_k = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^k - \left(\frac{1-\sqrt{5}}{2}\right)^k}{\sqrt{5}}$

and $f_{k+1} = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^{k+1} - \left(\frac{1-\sqrt{5}}{2}\right)^{k+1}}{\sqrt{5}}$

Induction case:

NTS: $f_{k+2} = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^{k+2} - \left(\frac{1-\sqrt{5}}{2}\right)^{k+2}}{\sqrt{5}}$

$$f_{k+2} = f_k + f_{k+1}$$

$$\therefore f_{k+2} = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^k - \left(\frac{1-\sqrt{5}}{2}\right)^k}{\sqrt{5}} + \frac{\left(\frac{1+\sqrt{5}}{2}\right)^{k+1} - \left(\frac{1-\sqrt{5}}{2}\right)^{k+1}}{\sqrt{5}}$$

(Proposition) Let n be a positive integer. Then $2^0 + 2^1 + \dots + 2^{n-1} = 2^n - 1$

Proof.

- Base case: $n=1 \Rightarrow 2^1 - 1 = 2^0 = 2^1 - 1 = 1$
- Induction hypothesis: Suppose $2^0 + 2^1 + \dots + 2^{k-1} = 2^k - 1$ (i)
- Induction case: NTS $2^0 + 2^1 + \dots + 2^{k-1} + 2^k = 2^{k+1} - 1$

Adding 2^k to both sides of (i):

$$2^0 + 2^1 + \dots + 2^{k-1} + 2^k = 2^k - 1 + 2^k = 2 \cdot 2^k - 1 = 2^{k+1} - 1 \quad \square$$

(Proposition) Let n be a natural number. Then $4^n - 1$ is divisible by 3.

(Proof) For each n , let $P(n)$ be the statement $3 \mid 4^n - 1$

- Base case:

- $P(0)$ is true because $4^0 - 1 = 0 = 0 \cdot 3$

- Inductive hypothesis:

- Suppose $P(k)$ is true $\forall k = n$:

- $3 \mid 4^k - 1 \therefore 4^k - 1 = 3m$ for some $m \in \mathbb{Z}$.

$$\therefore 4^k = 3m + 1 \quad (1)$$

- Inductive case:

- We must show that $4^{k+1} - 1$ is also divisible by 3.

- Multiply both sides of eq. (1) by 4:

$$4 \cdot 4^k = 4(3m + 1)$$

$$4^{k+1} = 12m + 4$$

- Subtract 1 from both sides:

$$4^{k+1} - 1 = 12m + 3$$

$$4^{k+1} - 1 = 3(4m + 1) = 3n \quad \text{where } n = 4m + 1$$

- Thus $P(k+1)$ is true.

- Therefore, $P(n)$ is true $\forall n \in \mathbb{N}$. \square

Chicken McNugget Theorem:

- If a and b are positive integers such that $\gcd(a, b) = 1$ (i.e., relatively prime)
- Then any number $n \geq ab - a - b + 1$ can be written as $ax + by$ with $(x, y) \in \mathbb{Z} : x \geq 0, y \geq 0$.

- Let $a = 4, b = 9$

$$- 24 = 6 \cdot 4$$

$$- 25 = 4 \cdot 4 + 1 \cdot 9$$

$$- 26 = 2 \cdot 4 + 2 \cdot 9$$

$$- 27 = 0 \cdot 4 + 3 \cdot 9$$

$$- 28 = 7 \cdot 4$$

$$- 29 = 5 \cdot 4 + 1 \cdot 9$$

$$- 30 = 3 \cdot 4 + 2 \cdot 9$$

$$- 31 = 1 \cdot 4 + 3 \cdot 9$$

$$- 32 = 8 \cdot 4$$

$$- 33 = 6 \cdot 4 + 1 \cdot 9$$

$$- 34 = 4 \cdot 4 + 2 \cdot 9$$

$$- 35 = 2 \cdot 4 + 3 \cdot 9$$

$$- 36 = 9 \cdot 4$$

largest number we can't write in the form $ax + by$ is $ab - a - b$

$$\text{If } n = 4a + 9b$$

$$\text{Then, } n + 4 = 4(a + 1) + 9b$$

5-step induction

$$n \rightarrow n + 5$$

Base cases:

$$s(0), s(1), \dots, s(4)$$

\therefore By 4-step induction, every number ≥ 24 can be written as $4a + 9b$ with $a, b \geq 0$.

- So, what's the largest amount of nuggets we can't buy with boxes of 4 and 9?

$$n = 4 \cdot 9 - 4 - 9 + 1$$

$$= 36 - 4 - 9 + 1 = 24$$

Strong Induction:

1. Base case:

- $S(n_0)$ is true

2. Strong Induction Hypothesis:

- Suppose $S(m)$ is true for all $n_0 \leq m \leq k$.

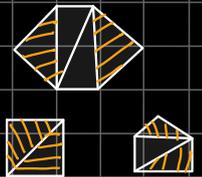
3. Induction step:

- Show $S(k+1)$ is true.

- Any polygon can be triangulated

- An external triangle is a triangle

in the triangulation that uses two sides.



- (Theorem): Any triangulation of a polygon with $n \geq 4$ sides has at least 2 external triangles.

- Strong induction hypothesis: Suppose any triangulation of an n -sided polygon has at least 2 ext. triangles.

- Induction case:

- Take a $k+1$ -sided polygon. Triangulated, the triangulation consists of drawing diagonals.

- Since $k+1 > 4$, it must have at least one diagonal.

- Isolate a diagonal d .



- Case 1:

- A or B is a triangle. The other one has at least 4 sides.

- The one with k sides by S.I.H has at least 2 external sides.

- Case 2: A and B have at least 4 sides

(Fundamental Theorem of Arithmetic):

- Any positive integer $n \geq 1$ can be factored as a product of primes in a unique way up to ordering.

- We'll prove that they can all be factored, we want prove factorization is unique.

Proof:

- $P(a)$ is true because a is itself a prime.

- Assume $P(k)$ is true $\forall k: 2 \leq k \leq n$.

- We need to prove $P(k+1)$, i.e., $k+1$ has a prime factorization

- Consider $k+1$, there are 2 cases:

1. If $k+1$ is prime:

- Then $k+1$ is already expressed as a product of prime factors (itself), and $P(k+1)$ holds.

2. If $k+1$ is composite:

- Then $\exists a, b \in \mathbb{Z}^+$ such $a \leq k, b \leq k$ and $k+1 = a \times b$

- By the inductive hypothesis, $P(a)$ and $P(b)$ are true.

- Therefore, $a = p_1 \times p_2 \times \dots \times p_m$ and $b = q_1 \times q_2 \times \dots \times q_n$ where p_i and q_j are prime.

- Hence $k+1 = a \times b = (p_1 \times p_2 \times \dots \times p_m)(q_1 \times q_2 \times \dots \times q_n)$

- This expresses $k+1$ as a product of prime factors, proving $P(k+1)$.

- By the principle of strong induction, $P(n)$ holds for all integers $n \geq 2$. \therefore every positive integer greater than 1 can be expressed as a product of prime factors.

20. Proof by Contradiction → if a claim is about something "not" being something else

- If we want to prove an "If A, then B" statement.
- Direct proof: Assume A is true, prove B is true.
- Contrapositive: Assume B is false, prove A is false.
- Contradiction: Assume A is true and B is false ($\neg B$ is true)

- ↓
- ① Assume $\neg P$
 - ② Find contradiction $P \wedge \neg P$
 - ③ Claim $\neg \neg P = P$
- Then you reach a contradiction.
"reductio ad absurdum"
which proves B had to be true.

(Theorem) There are infinitely many primes

Proof: Suppose, for the sake of contradiction, there are finitely many primes.

- Let p_1, p_2, \dots, p_n be the complete list of primes.
- Let $N = p_1 p_2 \dots p_n$
- $N > p_1, p_2, \dots, p_n$ so N is not prime because it can be factored as a product of primes.
- Let p be a prime divisor of N .
- Then $p = p_i$ for some $i = 1, 2, \dots, n$
- So $p \mid p_1 p_2 \dots p_n$ and $p \mid N = p_1 p_2 \dots p_n$
- So $p \mid N - p_1 p_2 \dots p_n = 1$!
- Thus, our assumption that there are infinitely many primes is false. \square

(Theorem) \sqrt{a} is irrational

Note: A number is rational if it can be written as a ratio (fraction) of two integers.

- Proof: - Assume, for contradiction, that $\sqrt{a} = \frac{p}{q}$ where p and q are integers in lowest terms.
- In particular, both p and q cannot be even.
 - Square both sides: $a = \frac{p^2}{q^2}$
 - Multiply by q^2 : $aq^2 = p^2$
 - p^2 is even, so p is even (since odd² = odd)
 - Let $p = 2k$
 - Substitute: $aq^2 = (2k)^2 = 4k^2$
 - q^2 is even, so q is even. !
 - Both p and q are even, contradicting that they're in lowest terms.
 - Thus, \sqrt{a} cannot be rational and must be irrational. \square

Prove that $(A-B) \cap (B-A) = \emptyset$

1. Assume, for contradiction, that $(A-B) \cap (B-A) \neq \emptyset$
- a. This means $\exists x$ such that $x \in (A-B) \cap (B-A)$
- b. So, $x \in A-B$ and $x \in B-A$
- c. $x \in (A-B)$ means $x \in A$ and $x \notin B$
- d. $x \in (B-A)$ means $x \in B$ and $x \notin A$
- e. From steps c and d, we have:
 - $x \in A$ and $x \notin A$!
 - $x \in B$ and $x \notin B$!
- f. Both statements in step e are contradiction.
- g. Therefore, our initial assumption must be false and we conclude that $(A-B) \cap (B-A) = \emptyset$.

(Theorem) No integer is both even and odd.

Proof:

1. Assume, for contradiction, that $\exists n \in \mathbb{Z}$ that is both even and odd.
2. This means $n = 2k$ and $n = 2m+1$ for some integers k, m .
3. Therefore, $2k = 2m+1$
4. Subtracting $2m$ from both sides:

$$2k - 2m = 1$$

$$2(k-m) = 1$$

$$k-m = \frac{1}{2}$$
5. Since $k-m$ is the subtraction of integers, $\frac{1}{2}$ is an integer. !
6. Therefore, our initial assumption must be false and we conclude that no integer is both even and odd.

(Theorem) e is irrational

Recall, $e = 1 + \frac{1}{1!} + \frac{1}{2!} + \dots = \sum_{k=0}^{\infty} \frac{1}{k!}$

Proof:

1. Assume, for contradiction, that $e = \frac{p}{q}$ where p and q are integers in lowest terms
- a. Multiply by $q!$:

$$e \cdot q! = \frac{pq!}{q}$$

$$e \cdot q! = p \cdot (q-1)!$$
3. Since p and q are integers, $e \cdot q!$ is an integer.
4. $e = 1 + \frac{1}{1!} + \frac{1}{2!} + \dots + \frac{1}{(q-1)!} + \frac{1}{q!} + \frac{1}{(q+1)!} + \dots$

$$e \cdot q! = \underbrace{q! + q! + \frac{q!}{2} + \frac{q!}{3} + \dots + \frac{q!}{q} + \frac{q!}{(q+1)} + \dots}$$

- $\sqrt{2}$ is a root of $x^2 - 2$
- $\sqrt[3]{2}$ is a root of $x^3 - 2$
- i is a root of $x^2 + 1$

when α is a root of a polynomial with integer coefficients, we say α is **algebraic**.

- If there is no polynomial with integer coefficients with root α , we say α is **transcendental**.

Examples: e , π , Liouville's number ^{first to be discovered (1840s)}

proved in 1750s

proved in 1800s (Hermite)

- proved that if α is algebraic, e^α is not algebraic

2.1. Smallest Counter-Example

- Steps:
- ① Negation: Assume the formula fails for some smallest k .
 - ② Base case: Verify the base case (typically $n=1$)
 - ③ Use the assumption for $n=k-1$
 - Since k is the smallest counterexample, the theorem must hold for all values $n < k$, including $n=k-1$.
 - ④ Express the quantity for $n=k$ using $n=k-1$.
 - ⑤ Simplify the expression and check if it contradicts the assumption that the theorem fails at k .
 - ⑥ Contradiction: Since we have shown that the proposition holds for $n=k$, this contradicts the assumption that the theorem fails at the smallest counterexample k . Therefore, no such smallest counterexample can exist.

Prove $1+2+\dots+n = \frac{n(n+1)}{2}$

Theorem: For all $n \in \mathbb{N}$:

$$S(n) = 1+2+\dots+n = \frac{n(n+1)}{2}$$

Proof:

- Assume $S(n)$ is false for some $n \in \mathbb{N}$. Let k be the smallest counterexample such that:

$$S(k) \neq 1+2+\dots+k = \frac{k(k+1)}{2} \quad (1)$$
 - $S(1)$ holds because $1 = \frac{1(1+1)}{2} = \frac{2}{2} = 1$, so $k > 1$.
 - Since k is the smallest counterexample, $S(k-1)$ holds such that:

$$S(k-1) = 1+2+\dots+(k-1) = \frac{(k-1)k}{2}$$
 - Notice that

$$S(k) = S(k-1) + k = \frac{(k-1)k}{2} + k = \frac{(k-1)k}{2} + \frac{2k}{2} = \frac{k(k-1+2)}{2} = \frac{k(k+1)}{2}$$
- contradicting (1) \square

(Well-Ordering Principle): If S is a non-empty subset of \mathbb{N} , then S has a smallest element.

(Theorem) Division Algorithm

Let a and b be integers

Then, there exist unique integers q, r such that

- ① $a = bq + r$
- ② $0 \leq r < b$

Proof:

- Suppose $b \nmid a$.
- Let $S = \{x \in \mathbb{N} \mid x = a - bq \text{ for } q \in \mathbb{Z}\}$
- S is non-empty because $a \in S$.
- $S \subseteq \mathbb{N}$.
- By WOP, there is a smallest element r in S .
- $r = a - bq$ so $a = bq + r$
- Assume, for contradiction, that $r \geq b$.
- Now, consider $r - b$.
- Since $b \nmid a$, $r - b \neq 0$ so $r - b > 0$ so $r - b \in \mathbb{N}$
- $r - b < r$, so $r - b \notin S$.
- If $r = a - bq \Rightarrow r - b = a - bq - b = a - b(q+1) = a - b(q_1)$
- So $r - b \in S$
- So $r < b$.

We proved existence, now let's prove uniqueness.

- Assume, for contradiction that $a = bq_1 + r_1$ with $0 \leq r_1 < b$ and $a = bq_2 + r_2$ with $0 \leq r_2 < b$

with $(q_1, r_1) \neq (q_2, r_2)$

- $bq_1 + r_1 = bq_2 + r_2$
- $b(q_1 - q_2) = r_2 - r_1$
- so $b \mid r_2 - r_1$
- Since $0 \leq r_1 < b$ and $0 \leq r_2 < b$,

$$-b < r_2 - r_1 < b$$
- so $r_2 - r_1 = 0$
- so $r_2 = r_1$

- Therefore, $bq_1 + r_1 = bq_2 + r_2$
- $bq_1 = bq_2$
- $\therefore q_1 = q_2$

(Euclidean Algorithm)

Let a, b be positive integers.

$$\log_a n \begin{cases} a = bq_1 + r_1 & 0 < r_1 < b \\ b = r_1q_2 + r_2 & 0 < r_2 < r_1 \\ \vdots \\ r_{n-2} = r_{n-1}q_{n-1} + r_n & 0 < r_n < r_{n-1} \\ r_{n-1} = r_nq_{n+1} + 0 \end{cases} \rightarrow \text{gcd}(a, b)$$

(Bezout's Identity): Given $a, b \in \mathbb{N}$, there exist $(x, y) \in \mathbb{Z}$ such that $ax + by = \text{gcd}(a, b)$

(Euclid's Lemma): If $p|ab$ where p is prime, $(a, b) \in \mathbb{N}$, then, $p|a$ or $p|b$.

(Proposition) Every natural number is even or odd.

$$P(n) = \forall n \in \mathbb{N}, n \text{ is even or odd.}$$

Proof:

- Assume, for contradiction, that $P(n)$ is false.
- Let k be the smallest counterexample such that:
 - $\exists k \in \mathbb{N}, k$ is neither even nor odd
 - i.e., $k \neq 2m$ or $k \neq 2m+1$ where $m \in \mathbb{Z}$.
- $P(1)$ holds because $1 = 2(0)+1$, so $k > 1$.
- Since k is the smallest counterexample, assume $P(k-1)$ holds s.t.:
 - $\exists k-1 \in \mathbb{N}, k-1$ is either odd or even.
- If $k-1$ is odd: $k-1 = 2m+1$
 - $k = 2m+2 = 2(m+1)$
 - so k is even ∇ (k is neither even or odd)
- If $k-1$ is even: $k-1 = 2m$
 - $k = 2m+1$
 - so k is odd ∇ (k is neither even or odd).
- Either way, we have a contradiction. Therefore $\forall n \in \mathbb{N}, n$ has to be even or odd. \square

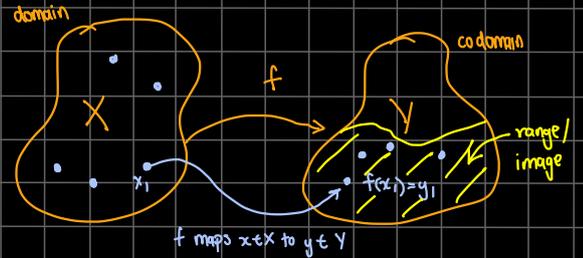
(Proposition) Every $n \geq 2$ is a product of one or more primes

(Proof):

- Assume, for contradiction, that $P(n)$ is false.
- Then, let k be the smallest counterexample s.t.:
 - $\exists k \geq 2$ that is not a product of primes.
 - so, $k = a \cdot b$ where $1 < a, b < k$
- $P(a)$ holds because a is itself a prime. so $k > a$.
- k is the smallest counterexample so $P(b)$ holds for $a \geq n > k$.
- Since $1 < a, b < k$, $P(a)$ and $P(b)$ hold. That is, a and b are a product of primes
- Since $k = a \cdot b$, k is a product of primes ∇
- Thus, $\forall n \geq 2, n$ is a product of primes. \square

2.4. Functions:

- $f: X \rightarrow Y$ (domain maps to codomain)
- f is a function from X to Y if
 - $\forall x \in X$ there exists a unique $y \in Y$ s.t. $f(x) = y$.
 - OR
 - A relation f is called a function provided $(a, b) \in f$ and $(a, c) \in f \Rightarrow b = c$



Examples:

- $f(x) = x^2$
 - Notice that this is a $f: \mathbb{R} \rightarrow \mathbb{R}^+$
 - $f(2) = 4 \nabla f(-2) = 4$
 - codomain = \mathbb{R}^+
 - Range: $\{x^2 \mid \sqrt{x}$ is an integer $\}$
 - $3 \notin \text{im } f$.

Useful Functions to Know:

Characteristic Function:

$$f(x) = \begin{cases} 0 & \text{if } x \notin A \\ 1 & \text{if } x \in A \end{cases} \text{ like truth table}$$

$$f: \mathbb{R} \Rightarrow \{0, 1\}$$

Identity Function:

$$\text{id}_A: A \Rightarrow A$$

where $\text{id}_A(a) = a \Rightarrow \text{example: } f(x) = x$

Floor & Ceiling Functions:

$$\text{Floor: } f(x) = \lfloor x \rfloor \rightarrow \lfloor 2.2y \rfloor = 2$$

$$\text{Ceiling: } f(x) = \lceil x \rceil \rightarrow \lceil 2.2y \rceil = 3$$

Injective, Surjective & Bijective Functions:

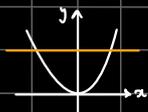
Injective: $f: X \rightarrow Y$ is injective (one-to-one)
 if $f(x_1) = f(x_2) \Rightarrow x_1 = x_2$ every value in the range corresponds to exactly one element in the domain.
 or if $x_1 \neq x_2$, then $f(x_1) \neq f(x_2)$

Examples: no two inputs can share an output.
 horizontal line test

1. $f: \mathbb{R} \rightarrow \mathbb{R}$
 $f(x) = 2x+1$
 $f(x) = f(y) \Rightarrow 2x+1 = 2y+1 \Rightarrow x = y \therefore f$ is 1-1.

2. Is $f(x) = x^2$ injective?

- Suppose $f(a) = f(b)$
 $a^2 = b^2$
 $\pm a = \pm b$



- Let $a = 2$ and $b = -2$.
 - So $f(a) = f(b) = 4$
 $f(2) = f(-2) = 4$.
 - Thus $f(a) = f(b)$ but $a \neq b \therefore f(x) = x^2$ is not injective. \square

3. Show $f(x) = 3x-2$ is injective.

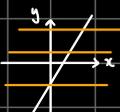
NTS $f(x_1) = f(x_2) \Rightarrow x_1 = x_2$

- Direct: Suppose $f(a) = f(b)$

$$3a-2 = 3b-2$$

$$3a = 3b$$

$$\therefore a = b \quad \square \text{ injective}$$



Surjective: $f: X \rightarrow Y$ is onto or surjective
 iff $\forall y \in Y, \exists x \in X$ s.t. $f(x) = y$.

im(f) = codomain of f range

every element in the codomain maps to at least one element in the domain.

- **Examples:**

1. $f: \mathbb{Z} \rightarrow \mathbb{Z}$, $f(x) = 2x+1$, is f onto?

- NTS $\forall y \in \mathbb{Z}, \exists x \in \mathbb{Z}$ s.t. $f(x) = y$
 - That is, $y = 2x+1$
 - Rearranging: $x = \frac{y-1}{2}$
 - $x \in \mathbb{Z} \forall x = 2k$ for some $k \in \mathbb{Z}$ (x has to be odd)
 - Since even integers are not in the image of f ,
 f is not onto (the codomain has to = the image)

2a. Show $f(x) = 5x+2$ is surjective when $f: \mathbb{R} \rightarrow \mathbb{R}$

- NTS $\forall y \in \mathbb{R}, \exists x \in \mathbb{R} : 5x+2 = y$
 - Solving for x : $x = \frac{y-2}{5}$
 - Since $y \in \mathbb{R}, x \in \mathbb{R}$ because subtraction & division
 b a non-zero constant (5) preserves real numbers. \square

b. What about when $f: \mathbb{Z} \rightarrow \mathbb{R}$?

- In this case, $x = \frac{y-2}{5} \notin \mathbb{Z} \forall y \in \mathbb{R}$
 - For example, if $y = 3, x = \frac{1}{5} \notin \mathbb{Z}$.
 - Thus $\exists y \in \mathbb{R}$ s.t. $\nexists x \in \mathbb{Z}$ with $f(x) = y$
 when $f: \mathbb{Z} \rightarrow \mathbb{R}$, f is not surjective.

c. What is the image of f ?

$$\text{Im}(f) = \{ y \in \mathbb{R} \mid y \equiv 2 \pmod{5} \}$$

d. What about when $f: \mathbb{R} \rightarrow \mathbb{Z}$?

- NTS $\forall y \in \mathbb{Z}, \exists x \in \mathbb{R}$ s.t. $f(x) = 5x+2$
 - $x = \frac{y-2}{5}$
 - Since $y \in \mathbb{Z}$ and subtraction, and division by a nonzero constant (5) yields a real number, $x \in \mathbb{R}$ and f is onto. \square

Bijective: A function is bijective if it is both one-to-one correspondence
 injective & surjective.

- For $f: X \rightarrow Y$, each $x \in X$ maps to exactly one unique $y \in Y$.
 - As a result $|X| = |Y|$

Function Proofs:

- $f: A \rightarrow B$

1. Is f 1-1?

① To prove: Assume $f(x) = f(y)$

Prove $x = y$

② To disprove: Find $x \neq y$ s.t. $f(x) = f(y)$

2. Is f onto? \rightarrow 1. let $f(x) = y$ 2. find $y = f^{-1}(x)$ 3. range of y 4. check range = codom

① To prove: Let $b \in B$

Find $a \in A$ s.t. $f(a) = b$.

② To disprove: Find $b \in B$ s.t. there is no $a \in A$ with $f(a) = b$.

Inverse Functions:

- Let f be a bijection from set A to B
- The inverse of f , f^{-1} , is the function from B to A defined as $f^{-1}(y) = x$ iff $f(x) = y$.
- In order for a function to be invertible, it must be a bijection.

(Theorem): Let $f: A \rightarrow B$

$$f^{-1}: \text{Im } f \rightarrow A$$

f^{-1} is a function iff f is 1-1.

(\Rightarrow):

- Let's prove if f is not 1-1, then f^{-1} is not a function.
- $\exists a_1, a_2 : a_1 \neq a_2$ s.t. $f(a_1) = f(a_2) = b$.
- So $(a_1, b), (a_2, b) \in f$.
- Thus, $(b, a_1), (b, a_2) \in f^{-1}$.
- Since $a_1 \neq a_2$, f^{-1} is not a function.

(\Leftarrow):

- Suppose f is 1-1. NTS: f^{-1} is a function.
- Suppose $(b, a_1), (b, a_2) \in f^{-1}$.
- (c.o.c.): show $a_1 = a_2$.
- Therefore $(a_1, b), (a_2, b) \in f$.
- $f(a_1) = b, f(a_2) = b$, so $f(a_1) = f(a_2)$.
- Since f is 1-1, $a_1 = a_2$.

(Theorem) If $f: A \rightarrow B$ and $g: B \rightarrow C$ are surjective (onto) functions, then $g \circ f$ is onto.

(Proof): $g \circ f: A \rightarrow C$

- Let $c \in C$. NTS $\exists a \in A$ s.t. $g(f(a)) = c$.
- Since g is onto, $\exists b \in B$ s.t. $g(b) = c$.
- Since f is onto, $\exists a \in A$ s.t. $f(a) = b$.
- So $g \circ f(a) = g(f(a)) = g(b) = c$. \square

(Theorem) $f: A \rightarrow B, g: B \rightarrow C, h: C \rightarrow D$

Then,

$$h \circ (g \circ f) = (h \circ g) \circ f$$

(Proof): NTS: 1. $\text{dom}(h \circ (g \circ f)) = \text{dom}((h \circ g) \circ f)$

$$a. h \circ (g \circ f)(a) = (h \circ g) \circ f(a)$$

$$1. \text{dom}(h \circ (g \circ f)) = \text{dom}(g \circ f) = \text{dom}(f)$$

$$\text{dom}((h \circ g) \circ f) = \text{dom}(f) \quad \checkmark$$

$$a. h \circ (g \circ f)(a) = h(g(f(a))) \quad \checkmark$$

$$(h \circ g) \circ f(a) = h \circ g(f(a))$$

$$= h(g(f(a))) \quad \checkmark$$

- Suppose f and f^{-1} are bijections, $f: A \rightarrow B, f^{-1}: B \rightarrow A$.

$$f \circ f^{-1} = \text{id}_B$$

$$f^{-1} \circ f = \text{id}_A$$

↳ if functions (i.e., f is not onto), then $f^{-1}: \text{Im}(f) \rightarrow A$
 $\therefore f \circ f^{-1} = \text{Im } f$

26. Composition:

- $f: A \rightarrow B$
 - $g: B \rightarrow C$
 - $g \circ f: A \rightarrow C$
- where $g \circ f(a) = g(f(a))$
- $g \circ f$ is defined if $\text{dom}(f) \subseteq \text{codom}(g)$ i.e., $(A \subseteq C)$

$$\left. \begin{array}{c} A \xrightarrow{f} B \xrightarrow{g} C \\ \searrow \quad \nearrow \\ \quad g \circ f \end{array} \right\} \begin{array}{l} \text{dom } g \circ f = \text{dom } f \\ \text{Im}(g \circ f) \subseteq \text{Im}(g) \end{array}$$

(Theorem) If $f: A \rightarrow B$ and $g: B \rightarrow C$ are injective (1-1) functions, then $g \circ f$ is injective.

(Proof): Suppose $g \circ f(a) = g \circ f(b)$. NTS $a = b$.

$$- g(f(a)) = g(f(b))$$

$$- \text{Since } g \text{ is 1-1, } f(a) = f(b) \quad \leftarrow \text{we can cancel out the functions, leaving the inputs}$$

$$- \text{Since } f \text{ is 1-1, } a = b. \quad \square$$

Counting Functions:

- $A = \{1, 2, \dots, n\}$
- $B = \{1, 2, \dots, m\}$

in a function, every element in the domain maps to something in the codomain.

1. How many functions $f: A \rightarrow B$ exist?

- $f(1)$ has m possibilities
- $f(2)$ has m "
- $f(n)$ has m "



$$m^n$$

2. How many 1-1 functions $f: A \rightarrow B$?

- $f(1)$ has m
- $f(2)$ has $m-1$
- \vdots
- $f(n)$ has $(m-n)+1$

$\binom{m}{n} \cdot n!$ if $m \geq n$
 0 if $m < n$
 map n elements onto m elements
 ways of arranging
 we don't have enough elements in the domain to map to every element in the codomain.

3. How many onto functions $f: A \rightarrow B$?

inclusion-exclusion

all functions - functions that aren't onto

$$m^n - \binom{m}{1} (m-1)^n + \binom{m}{2} (m-2)^n - \binom{m}{3} (m-3)^n + \dots + (-1)^{m-1} \binom{m}{m-1} (m-1)^n$$

each missing value
 missing a value
 pair
 triple

$$\sum_{k=0}^{m-1} (-1)^k \binom{m}{k} (m-k)^n = n! \text{ when } m=n$$

a. Prove (1) + (3) \Rightarrow (2):

- let $B = \{b_1, b_2, \dots, b_n\}$, so $|B| = n$
- Since f is onto for each b_i , there exists $a_i \in A$ such that $f(a_i) = b_i$.
- From the definition of a function, a_1, a_2, \dots, a_n are all distinct.
- $\{a_1, a_2, \dots, a_n\} \subseteq A$
- $|\{a_1, a_2, \dots, a_n\}| = n = |A|$
- So, $A = \{a_1, a_2, \dots, a_n\}$.
- Suppose $f(a_i) = f(a_j)$ so $x = a_i$ for some i
 $y = a_j$ for some j .
- $f(a_i) = b_i, f(a_j) = b_j$ & $b_i = b_j \Rightarrow i=j = a_i = a_j$ so f is 1-1.

b. Prove (2) + (3) \Rightarrow (1):

- let $A = \{a_1, a_2, \dots, a_n\}$.
- let $b_1 = f(a_1) \in B$
- $b_2 = f(a_2) \in B$
- \vdots
- $b_n = f(a_n) \in B$
- Since f is 1-1, b_1, b_2, \dots, b_n are all distinct.
- We know f is onto, so there is no other output.
- so $B = \{b_1, b_2, \dots, b_n\}$
- $|A| = n = |B|$

- (Theorem) Let A, B be finite sets and $f: A \rightarrow B$.
- If 2 of the following 3 are true, the third is also true.
 - 1) $|A| = |B|$
 - 2) f is one-to-one
 - 3) f is onto

1. Prove (1) + (2) \Rightarrow (3):

- let $A = \{a_1, a_2, \dots, a_n\}$
- $f(a_1) = b_1$
- $f(a_2) = b_2$
- \vdots
- $f(a_n) = b_n$
- Since f is 1-1, b_1, b_2, \dots, b_n are all distinct
- $\{b_1, b_2, \dots, b_n\} \subseteq B$.
- But $|B| = |A| = n$
 and $|\{b_1, b_2, \dots, b_n\}| = n$
- so $B = \{b_1, b_2, \dots, b_n\}$.
- Take $b \in B$ so $b = b_i$ for some i .
- Then $f(a_i) = b_i$, so f is onto.

2.5. Pigeonhole Principle:

if A & B are finite sets and $|A| > |B|$, there can be no 1-1 functions $f: A \rightarrow B$

- If you have $n+1$ pigeons flying towards n pigeonholes, at least 2 pigeons will share a pigeonhole.

Generalized Pigeonhole Principle:

- Given n items and k boxes, if $m > n$, there is at least one box with $\lceil \frac{n}{k} \rceil$ items.

Examples:

1. 28 students, 12 months

$$\lceil \frac{28}{12} \rceil = \lceil 2.33 \rceil = 3$$

\therefore At least 3 people have a birthday in 1 month.

2. 50,000 people, 366 possible birthdays

$$\lceil \frac{50,000}{366} \rceil = \lceil 136.6 \rceil = 137$$

\therefore At least 137 people share a birthday.

3. Show that at a party at least 2 people have the same number of acquaintances.

- Ex: 7 people

0, 1, 2, 3, 4, 5, 6 } 7 options
 Know none Know everyone \Rightarrow can't both exist
 \therefore there are actually only 6 options.

4. Given 6 people at a party, show you can find either

- a) 3 of them that all know each other,
- b) 3 of them that don't know each other.



Ramsey Number $r(3,3) = 6$

Proof:

1. First, let's model this:

- Let's represent people as vertices.
- An edge between two people means they know each other, no edge means they don't.
- So, we're looking for either: a) A triangle (3 mutual edges) or b) an anti-triangle (3 vertices with no edges between them).

2. Let's focus on one person, call them A.

- Looking at the other 5 people, each of these 5 must either know or not know A.
- By pigeonhole, at least $\lceil \frac{5}{2} \rceil = 3$ people either know or don't know A.

3. Let's consider the cases:

- Case 1: 3 know A (call them B, C, D)

- Now look at edges between B, C, and D.

- If any 2 of them know each other, we get a triangle ABC or ABD or ACD.

- If none of them do, we get an anti-triangle BCD.

- Case 2: 3 don't know A (call them B, C, D)

- If any 2 don't know each other, we

get an anti-triangle ABC or ACD.

- If all 3 know each other, we get triangle BCD.

4. In either case:

- We find 3 people who all know each other or 3 people who don't. \square



5. Suppose we have 5 points in the plane with integer coordinates

- Show there is a point with integer coordinates in at least one segment connecting the original 5 points.

- I'll show that the midpoint of at least one pair has integer coordinates.

$$m = \left(\frac{x_1+x_2}{2}, \frac{y_1+y_2}{2} \right)$$

- For $m \in \mathbb{Z}$, x_1, x_2 and y_1, y_2 need to be even. Thus, x_1 and x_2 need to have the same parity. Same for y_1 and y_2 .

- For each midpoint (x_1, x_2, y_1, y_2) can be (even, even), (even, odd), (odd, odd), or (odd, even). So 4 possibilities.

- From the pigeonhole principle, there are at least $\lceil \frac{5}{4} \rceil = 2$ coordinates with the same parity.

- b. Rational numbers are numbers of the form $\frac{a}{b}$, $a, b \in \mathbb{Z}$, $b \neq 0$.

- Prove that every rational number must have a repeating decimal representation.

Proof:

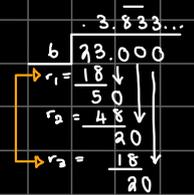
- Given a rational number $\frac{a}{b}$, when we perform long division, we multiply the remainder and divide by b . Each step gives us the next digit and a new remainder.

- When dividing by b , the possible remainders are $0, 1, a, \dots, b-1$. So there are only b possibilities.

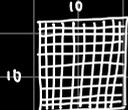
- After $b+1$ steps of long division, we have $b+1$ remainders but since there are only b possibilities, some remainder must repeat (by pigeonhole principle).

- When a remainder repeats, the next digits will be the same as what came after the first occurrence.

- Thus, every rational number must have a repeating decimal expansion. \square



7. Given a 10×10 chessboard and 41 rooks, prove you can find 5 rooks that are not attacking each other.



Proof:

- Since we have 41 rooks and 10 rows, one row has at least $\lceil \frac{41}{10} \rceil = 5$ rooks. This is the row with the most rooks.
- Let's call this row R_5 .

- Each of the 5 rooks are in a different column.

- Mark the C_1, C_2, C_3, C_4, C_5 . R_5

- Now we remove R_5 and consider the sub-board:

- So now we have a 9×5 board with 36 rooks.

- 45 spaces, 36 rooks, so 9 squares must be empty. By pigeonhole, 1 column must have at least $\lceil \frac{9}{5} \rceil = 2$ empty spaces.



① - Choose 1 of these empty spaces and remove the rook. *empty

- Now we have a 8×5 , so 40 spaces, 34 rooks

- so there are 8 empty spaces. So one row has at least

$\lceil \frac{8}{5} \rceil = 2$ empty spaces, choose 1. ②

- $7 \times 5 \Rightarrow 35$ spaces, 28 rooks, 7 empty. So, 1 row has at least

$\lceil \frac{7}{5} \rceil = 2$ empty, choose 1. ③

- $6 \times 5 \Rightarrow 30$ spaces, 24 rooks, 6 empty. So, 1 row has at least

$\lceil \frac{6}{5} \rceil = 2$ empty, choose 1. ④

- $5 \times 5 \Rightarrow 25$ spaces, 20 rooks, 5 empty. So,

1 row has at least $\lceil \frac{5}{5} \rceil = 1$ empty, choose it ⑤

- We've chosen 5 empty spaces in different rows and columns, essentially selecting 5 non-attacking rooks \square .

8. Given a 41×25 chessboard with an integer written on each square, show there is a rectangle (with width, length ≥ 2) in the board, whose 4 corner numbers sum to a multiple of 4.

Proof:

- we will color a square black if it's a multiple of 4.

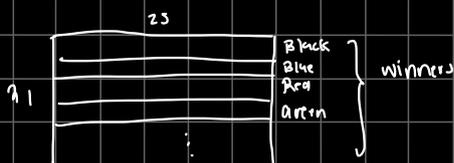
Blue if it's $1 \pmod 4$, Red if $2 \pmod 4$,

Green if it's $3 \pmod 4$.

- If we find a rectangle with corners of the same color, when we add them up, we will get a multiple of 4.

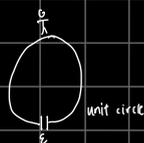


- We have 4 colors, so by pigeonhole, 1 color appears at least 7 times.



- There is a color that appears at least 7 times in the most rows. By pigeonhole, it will be the winner in at least 6 rows.

9. step size = α , $\alpha \notin \mathbb{Q}$
hole size $\varepsilon > 0$



- Prove that the person will fall through eventually.

- Proof:

- We need to find i and j such that

$$\lceil (j-i)\alpha \rceil < \varepsilon$$

$$\rightarrow \{n\alpha\} = 0, \alpha$$

- $\varepsilon > 0$, $\lim_{n \rightarrow \infty} \frac{1}{n} = 0$ so $\exists N \in \mathbb{N}$ s.t. $\frac{1}{N} < \varepsilon$.



- N arcs of the circle.

- If you have $N+1$ steps, 2 of them land on the same arc.

- Since $\alpha \notin \mathbb{Q}$, two step sizes can't fall on the same point: $i\alpha = N_i + \{i\alpha\}$

$$j\alpha = N_j + \{j\alpha\}$$

$$(j-i)\alpha = N_j - N_i$$

$$\therefore \alpha = \frac{N_j - N_i}{j - i} \in \mathbb{Q}$$

10. Let $S = \{1, 2, \dots, 20\}$

- If we pick 11 numbers, we are guaranteed that the sum of two picked numbers is 21.

- Pairs that equal 21:

$$\left. \begin{aligned} (1, 20), (2, 19), (3, 18), (4, 17), (5, 16), \\ (6, 15), (7, 14), (8, 13), (9, 12), (10, 11) \end{aligned} \right\} 10 \text{ pairs}$$

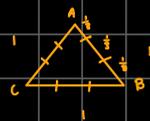
- Suppose you pick the first number in each pair such that none of the numbers you choose sum up to 21:

$$\{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$$

- You've picked 10 numbers, so you need to pick 1 more.

No matter which number you pick, you'll complete a pair that sums to 21. \square

11. Let $\triangle ABC$ be an equilateral triangle with $AB=1$. Show that by selecting 10 points on the perimeter, there are at least two with distance $\leq \frac{1}{3}$ apart.



Proof:

- If we divide the perimeter into sections of length $\frac{1}{3}$ we get 9 sections (including vertices A, B, C).

- By pigeonhole, at least $\lceil \frac{10}{9} \rceil = 2$ points must be in the same section.

- For two points in the same section, their arc length distance $\leq \frac{1}{3}$. \square

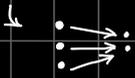
(Erdős-Szekeres): Let n be a positive integer. Every sequence of n^2+1 distinct integers must contain a monotone subsequence of length $n+1$.

Cardinality:

The cardinality of a set A , denoted by $|A|$ is the number of elements in A when A is a finite set.

For infinite sets:

- Let A and B be sets.
- $|A| = |B|$ if there exists a bijection $f: A \rightarrow B$
- $|A| \leq |B|$ if there exists a 1-1 function $f: A \rightarrow B$.
- $|A| \geq |B|$ if there exists an onto function $f: A \rightarrow B$.



3. Show $|C(0,1)| = |C(-1,1)|$

$f: (0,1) \rightarrow (-1,1)$

$f(x) = 2x-1$

- Injective: Suppose $f(a) = f(b)$

$2a-1 = 2b-1$

$2a = 2b \Rightarrow a = b \quad \square$

- Surjective: Let $y = f(x)$

$y = 2x-1$

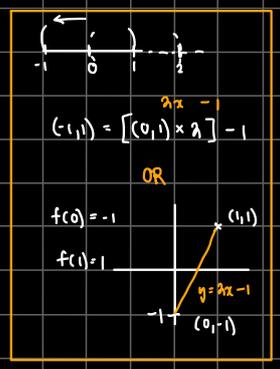
$x = \frac{y+1}{2} \quad f\left(\frac{y+1}{2}\right) = \frac{2(y+1)}{2} - 1 = y$

- NTS $\frac{y+1}{2} \in (0,1)$:

$0 < \frac{y+1}{2} < 1$

$0 < y+1 < 2$

$\therefore -1 < y < 1 \quad \square$



1. Let's show $|2\mathbb{Z}| = |\mathbb{Z}|$:

$f: \mathbb{Z} \rightarrow 2\mathbb{Z}$

$f(n) = 2n$

injective:

$f(a) = f(b)$

$2a = 2b$

$a = b$

surjective:

Let $y = f(x)$

$x = \frac{y}{2}$

$f(x) = \frac{2y}{2} = y \quad \square$

2. Show $|\mathbb{N}| = |\mathbb{Z}|$



$f: \mathbb{N} \rightarrow \mathbb{Z}$

$1 \rightarrow 0$

$2 \rightarrow 1$

$3 \rightarrow -1$

$f(n) = \begin{cases} \frac{n}{2} & \text{if } n \text{ is even} \\ \frac{1-n}{2} & \text{if } n \text{ is odd} \end{cases}$

- Prove $f(n)$ is bijective:

- Injective: Suppose $f(a) = f(b)$

- Case 1: both odd

$\frac{1-a}{2} = \frac{1-b}{2}$

$1-a = 1-b$

$a = b$

- Case 2: both even:

$\frac{a}{2} = \frac{b}{2}$

$a = b$

- Case 3: different parity

$\frac{a}{2} = \frac{1-b}{2}$

$a = 1-b$

$a+b = 1 \Rightarrow \Leftrightarrow a, b \in \mathbb{N}$ so $a+b \geq 1+1 = 2$

- So $f(a) \neq f(b)$

- Therefore, f is injective. \square

- Surjectivity: Let $y = f(x)$

- Case 1: x is even

$y = \frac{x}{2}$

$x = 2y$

$f(2y) = \frac{2y}{2} = y$

- Therefore, f is surjective. \square

- Case 2: x is odd

$y = \frac{1-x}{2}$

$x = 1-2y$

$f(1-2y) = \frac{1-(1-2y)}{2} = y$

(Theorem) Two open intervals have the same cardinality

i.e., $| (a,b) | = | (c,d) |$

$f(a) = c \quad \left\{ \begin{array}{l} (a,c) \longleftrightarrow (b,d) \\ m = \frac{d-c}{b-a} \end{array} \right.$

(Proof): $f(x) = \left(\frac{d-c}{b-a}\right)x + \left[c - \left(\frac{d-c}{b-a}\right)a\right]$

$f(a) = \left(\frac{d-c}{b-a}\right)a + c - \left(\frac{d-c}{b-a}\right)a = c$

$f(b) = \left(\frac{d-c}{b-a}\right)b + c - \left(\frac{d-c}{b-a}\right)a = c + \frac{d-c}{b-a}(b-a)$

Injective: Suppose $f(r) = f(s)$

$\frac{d-c}{b-a} \cdot r + c - \frac{d-c}{b-a} a = \frac{d-c}{b-a} \cdot s + c - \frac{d-c}{b-a} a$

$\frac{d-c}{b-a} \cdot r = \frac{d-c}{b-a} \cdot s$

- Since $b-a > 0$ and $d-c > 0$, we can multiply by both sides.

- Thus: $r = s \quad \square$

(Theorem) Equality of cardinality is an equivalence relation.

(Proof) We need to prove the relation is reflexive, symmetric, and transitive.

- Reflexive: Let A be a set. NTS $|A| = |A|$

- Take $f: A \rightarrow A$ defined by $f(a) = a$.

- $f(a) = f(b) : a = b \quad \square$ injective

- Let $y = x \quad \square$ surjective.

- Symmetric: Suppose $|A| = |B|$, NTS $|B| = |A|$

- $f: A \rightarrow B$ is a bijection. NTF $g: B \rightarrow A$.

- Take $g: B \rightarrow A$ defined by $g = f^{-1}$.

- Since f is a bijection, f^{-1} is also a bijection

- So, $|B| = |A| \quad \square$

- Transitive: Suppose $|A|=|B|$ and $|B|=|C|$, NTS $|A|=|C|$
 - There is a bijection $f: A \rightarrow B$.
 - There is a bijection $g: B \rightarrow C$.
 - Then, $g \circ f: A \rightarrow C$ is a bijection. \square



2. Show that $\{1, \frac{1}{4}, \frac{1}{9}, \frac{1}{16}, \dots\}$ is countable.
- Let $S = \{1, \frac{1}{4}, \frac{1}{9}, \frac{1}{16}, \dots\}$
- Let $f: \mathbb{N} \rightarrow S$
- $\therefore f(n) = \frac{1}{n^2}$

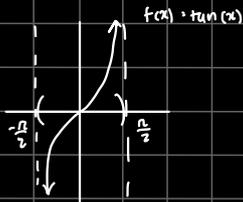
Injectivity:

- Suppose $f(a) = f(b)$
- $\frac{1}{a^2} = \frac{1}{b^2}$
- $a^2 = b^2 \Rightarrow a = b \therefore f$ is injective. \square

Surjectivity:

- Let $y = f(x)$
- $\therefore y = \frac{1}{x^2}$
- $x = \pm \frac{1}{\sqrt{y}}$
- $f(\pm \frac{1}{\sqrt{y}}) = y \therefore f$ is surjective.

4. Show $|(-\frac{\pi}{2}, \frac{\pi}{2})| = |\mathbb{R}|$



- $f: (-\frac{\pi}{2}, \frac{\pi}{2}) \rightarrow \mathbb{R}$
- defined by $f(x) = \tan(x)$

Injective:

- Suppose $\tan(a) = \tan(b)$ for some $a, b \in (-\frac{\pi}{2}, \frac{\pi}{2})$.
- For contradiction, assume $a \neq b$, and $a < b$.
- $\tan(x)$ is cont. on $[a, b]$ and diff. on (a, b) .
- By Rolle's Theorem, $\exists c \in (a, b)$

Rolle's Theorem:

- If f is cont. on $[a, b]$ and f is diff on (a, b) and $f(a) = f(b)$, $\exists c \in (a, b)$ s.t. $f'(c) = 0$.

- such that $f'(c) = 0$.
- $f'(c) = \sec^2(c) = 0$
- $\tan(x)$ is increasing on $(-\frac{\pi}{2}, \frac{\pi}{2})$
- Since our assumption that $a \neq b$ led to a contradiction, we conclude $a = b$.

- Thus, f is injective.

Surjective:

- Let $y = f(x) \therefore y = \tan x$
- $x = \tan^{-1}(y) \therefore f(\tan^{-1}(y)) = \tan(\tan^{-1}(y)) = y \quad \square$



(Theorem) If A and B are countable, then $A \cup B$ is countable

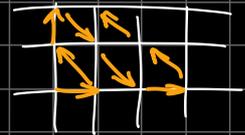
(Proof):

- Since A and B are countable, we can write $A = \{a_1, a_2, \dots\}$ and $B = \{b_1, b_2, \dots\}$
- Thus, $A \cup B = \{a_1, b_1, a_2, b_2, \dots\}$
- If $A \cap B \neq \emptyset$, don't list duplicates.

(Theorem) If A and B are countable, $A \times B$ is countable.

(Proof):

- $A = \{a_1, a_2, \dots\}$
- $B = \{b_1, b_2, \dots\}$
- $A \times B = \{(a_1, b_1), (a_2, b_1), \dots\}$



Countable:

smallest infinity

- A set A is countable if $|A| = |\mathbb{N}| = \aleph_0$ [aleph not]
- A set is countable if there is a bijection between that set and the natural numbers.

1. Show that $\{0, 2, 4, 6, 8, \dots\}$ is countable.

NTS: $\{0, 2, 4, 6, 8, \dots\} \leftrightarrow \{0, 1, 2, 3, \dots\}$

$$\begin{aligned} 2n &\mapsto n \\ n &\mapsto \frac{n}{2} \\ \dots f(n) &= \frac{n}{2} \end{aligned}$$

Verifying Bijectivity:

1. Injectivity:

- Suppose $f(a) = f(b)$
- $\frac{a}{2} = \frac{b}{2} \therefore a = b \therefore f$ is injective. \square

2. Surjectivity:

- Let $y = f(x)$
- $y = \frac{x}{2} \therefore x = 2y$
- $f(2y) = \frac{2y}{2} = y \therefore f$ is surjective. \square

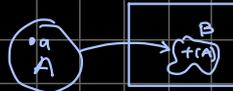
(Theorem) $|\mathbb{R}^{\mathbb{N}}| > |\mathbb{N}| \leftarrow$ Cantor's Theorem

(Proof):

$$\begin{aligned} |\mathbb{R}^{\mathbb{N}}| &= |\mathbb{R}| \\ \text{so } |\mathbb{R}| &\neq |\mathbb{N}| \\ &\downarrow \\ &\mathbb{R} \text{ is uncountable.} \end{aligned}$$

Cantor's Theorem:

- Recall, $|A| \leq |B|$ if there exists $f: A \rightarrow B$ that is one-to-one. Equivalently, if there exists $g: B \rightarrow A$ that is onto.



- $f: A \rightarrow B$ is 1-1

$$g: B \rightarrow A \quad g(x) = \begin{cases} f^{-1}(x) & \text{if } x \in \text{im}(f) \\ a & \text{if } x \notin \text{im}(f) \end{cases}$$

- $|A| < |B|$ if $|A| \leq |B|$ and $|A| \neq |B|$

(Cantor's Theorem): $|A| < |2^A|$

Rephrase: there exist no onto functions from A to 2^A .

To show $|A| < |2^A|$ is easy:

$$f: A \rightarrow 2^A \\ f(a) = \{a\}$$

Example: $A = \{1, 2, 3\}$

$$f(1) = \{1\} \quad 2^A = \{\emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$$

$$f(2) = \{2\}$$

$$f(3) = \{3\}$$

Prove $|A| \neq |2^A|$:

- Assume, for contradiction, that they have the same size.

$$|A| = |2^A|$$

- Let $f: A \rightarrow 2^A$.

- Let's show that f is not onto. [let's create a set $B \notin \text{im}(f)$]

- Let $B = \{x \in A \mid x \notin f(x)\}$

- (claim: $B \notin \text{im}(f)$)

- Suppose, for contradiction, that $B \in \text{im}(f)$.

- So, $\exists a \in A$ s.t. $f(a) = B$

- Is $a \in B$?

- If $a \in B$, then $a \notin f(a) = B$ \uparrow

- If $a \notin B$, then $a \in f(a)$, so $a \in B$. \uparrow

- Thus, $|A| < |2^A| < |2^{2^A}| < \dots$

Cantor-Bernstein-Schroeder Theorem:

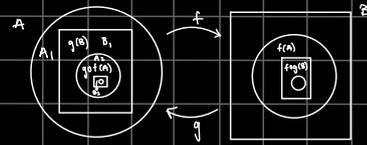


(Theorem) If $|A| \leq |B|$ and $|B| \leq |A|$, then $|A| = |B|$

(Proof): We have $f: A \rightarrow B$ that is 1-1.

We have $g: B \rightarrow A$ that is 1-1.

Goal: Find a bijection $h: A \rightarrow B$.



$$\left. \begin{aligned} - A_1 &= A - g(B) \\ - B_1 &= g(B) - g \circ f(A) \\ - A_2 &= g \circ f(A) - g \circ f \circ g(B) \\ - B_2 &= g \circ f \circ g(B) - g \circ f \circ g \circ f(A) \end{aligned} \right\} \begin{aligned} - A_{n+1} &= g \circ f \circ A_n \\ - B_{n+1} &= g \circ f \circ B_n \end{aligned}$$

- Let $C = A_1 \cup A_2 \cup \dots$

$$D = A - C$$

$$h(x) = \begin{cases} f(x) & \text{if } x \in C \\ g^{-1}(x) & \text{if } x \in D \end{cases} \quad h: A \rightarrow B$$

- Claim: h is a bijection

- Injectivity: Suppose $h(a) = h(b)$

- case 1: $a, b \in C$

$$h(a) = h(b) \Rightarrow f(a) = f(b)$$

$$f \text{ is 1-1} \Rightarrow a = b$$

- case 2: $a, b \in D$

$$h(a) = h(b) \Rightarrow g^{-1}(a) = g^{-1}(b)$$

$$g \text{ is 1-1} \Rightarrow a = b$$

- case 3: $a \in C, b \in D$

$$h(a) = h(b) \Rightarrow f(a) = g^{-1}(b)$$

$$a \in C, \text{ so } a \in A_i \text{ for some } i \quad \text{apply } g \text{ to both sides}$$

$$\Rightarrow g \circ f(a) \in A_{i+1}$$

$$g \circ f(a) = b \in A_{i+1} \Rightarrow b \in C \quad \uparrow$$

- Surjectivity: NTS $\forall b \in B, \exists a \in A$.

- Let $b \in B$.

- Let $E = f(A) \cup f \circ g(A) \cup \dots$

$$F = B - E$$

- If $b \in E$:

$$b \in f(A_i) \text{ for some } i, \text{ so } \exists a \in A_i \text{ s.t. } f(a_i) = b$$

$$\text{So, } h(a_i) = f(a_i) = b$$

- If $b \in F$:

$$b \notin f(A_i) \text{ for any } i$$

$$\text{So, } g(b) \notin g \circ f(A_i) \text{ so } g(b) \notin A_{i+1}$$

$$A_1 = A - g(B) \text{ so } g(b) \notin A_1$$

$$\text{So } g(b) \in A_1 \cup A_2 \cup \dots \Rightarrow g(b) \in D$$

$$\text{Thus, } h(g(b)) = g^{-1}(g(b)) = b, \text{ so } b \in \text{im}(h)$$

- Therefore, h is onto \square

(Corollary) $|C_0, I| = |C_0, I|$

- We showed that $|C_0, I| = |\mathbb{R}|$.

Proof:

- We know $|C_0, I| \leq |\mathbb{R}|$ since $f(x) = x$ where $f: [0, 1] \rightarrow \mathbb{R}$
- Since $|C_0, I| = |\mathbb{R}|$, $|\mathbb{R}| \leq |C_0, I|$
 $\therefore |C_0, I| = |\mathbb{R}|$.

(Theorem) $|\mathbb{R}| = 2^{\aleph}$

Proof: We'll show $|C_0, I| = 2^{\aleph}$

- Let $f: 2^{\aleph} \rightarrow [0, 1]$
 $f(A) = \sum_{a \in A} \frac{1}{2^a}$ $f(\emptyset) = 0$
 $f(N) = \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots = 1$

- Claim: f is onto.

- Let $x \in [0, 1]$

- WTS: $\exists A \in 2^{\aleph}$ s.t. $f(A) = x$.

$x = 0.b_1b_2b_3\dots$

Example: $x = 0.001001001001\dots_2$
 $= \frac{1}{2^3} + \frac{1}{2^6} + \frac{1}{2^9} + \frac{1}{2^{12}} + \dots$

$A = \{3, 6, 9, 12, \dots\}$

$f(A) = \frac{1}{2^3} + \frac{1}{2^6} + \frac{1}{2^9} + \frac{1}{2^{12}} + \dots = x$

- So f is onto.

- Thus, $2^{\aleph} \geq |C_0, I|$

- Now, consider $g: 2^{\aleph} \rightarrow [0, 1]$
 $g(A) = \sum_{a \in A} \frac{1}{10^a}$

$g(\{1, 2\}) = 0.11$

$g(\{3, 4, 6\}) = 0.001101$

- Thus, g is 1-1

- So, $2^{\aleph} \leq |C_0, I|$

- Since $2^{\aleph} \geq |C_0, I|$ and $2^{\aleph} \leq |C_0, I|$,

$2^{\aleph} = |C_0, I|$

Using the Cantor-Bernstein-Schroeder Theorem, prove $2^{\aleph} = |\mathbb{R}|$.

- $f: 2^{\aleph} \rightarrow [0, 1]$

$f(A) = \sum_{a \in A} \frac{1}{2^a}$ (e.g., $f(\{2, 5, 6\}) = 0.010011$)

- $g: [0, 1] \rightarrow 2^{\aleph}$

$g(x) = \sum_{a \in A} \frac{1}{10^a}$ $0.010110011\dots$
 $= 0.1011001$

- $A_1 = A - g(A)$ these are sets that end with $\{2, 5, 6, \dots\}$

- $A_2 = g \circ f(A_1)$

54. Fundamentals of Partially Ordered Sets

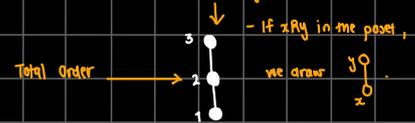
- A partially ordered set is a pair $P = (X, R)$ where X is a nonempty set and R is a relation on X that satisfies the following conditions:
 - R is reflexive: $\forall x \in X, x R x$.
 - R is antisymmetric: $\forall x, y \in X, x R y$ and $y R x \Rightarrow x = y$.
 - R is transitive: $\forall x, y, z \in X$, if $x R y$ and $y R z$, then $x R z$.

1. $X = \{\emptyset, \{1, 2, 3\}\}$, $R = \{(x, y) \mid x \subseteq y\}$

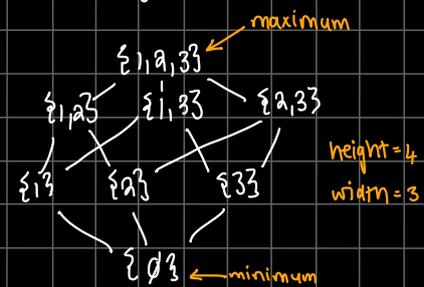
$R = \{(\emptyset, \emptyset), (\emptyset, \{1, 2\}), (\emptyset, \{1, 3\}), (\emptyset, \{2, 3\}), (\emptyset, \{1, 2, 3\}), (\{1, 2\}, \{1, 2, 3\}), (\{1, 3\}, \{1, 2, 3\}), (\{2, 3\}, \{1, 2, 3\})\}$



- Since R is reflexive, we don't have to draw the loops.
- Since R is transitive, we don't have to draw $1 R 3$.
- The result is a Hasse Diagram:



2. $X = 2^{\{1, 2, 3\}} = \{\emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$
 $R = \{(x, y) \mid x \subseteq y\}$



Definitions:

- Suppose $P = (X, \leq)$ is a poset.
 - comparable \Leftrightarrow
 - incomparable \nLeftrightarrow
 - below \leq
 - above \geq

(Incomparable): $a, b \in X$ are called incomparable if $a \not\leq b$ and $b \not\leq a$. $a \not\leq b$

(Chain): $C \subseteq X$ is a chain if for any $a, b \in C$, a, b are comparable.

(Anti-chain): $A \subseteq X$ is an antichain if for any $a, b \in A$, a, b are incomparable.

(Height): the height of P is the maximum size of a chain.
 - It is also the longest path from the minimums to the maximums.

(Width): the width of P is the maximum size of an antichain.
 - width = max (longest antichain at each level)

55. Max and Min

(Maximum): An element $m \in X$ is a maximum if for every $x \in X, x \leq m$.

(Minimum): An element $m \in X$ is a minimum if for every $x \in X, m \leq x$.

(Maximal): An element $m \in X$ is maximal if for every x that is comparable to $m, x \leq m$.

no successor

(Minimal): An element $m \in X$ is minimal if for every x that is comparable to $m, m \leq x$.

no predecessor

$$1. \quad X = \{1, 2, \dots, 50\}$$

$$R = \{(x, y) : x | y\}$$

maximals = 20, 50 \rightarrow for any number $x, m > x$ if

minimum = 1 $m \equiv 2^k$

height = 6

width = 25

$$2. \quad X = \{1, 2, \dots, n\}$$

$$R = \{(x, y) : x | y\}$$

height = $\lfloor \log_2 n \rfloor + 1$

width = $\lceil \frac{n}{2} \rceil$

maximals = $\lfloor \frac{n}{2} \rfloor - 1, \lfloor \frac{n}{2} \rfloor + 1, \dots, \lfloor \frac{n}{2} \rfloor + \lceil \frac{n}{2} \rceil$

minimum = 1

accounts for the 1

n

(Theorem): Given a set with $n+1$ positive integers $\leq 2n$, at least two of them divide each other.

(Proof):

- Every positive integer can be written as:

$$m = 2^k \cdot b \quad \text{where } b \text{ is odd.}$$

- How many odd numbers are there between 1 and $2n$?

$$\frac{2n}{2} = n$$

- By the pigeonhole principle, at least 2 numbers have the same odd part.

$$2^i \cdot b \quad \text{and} \quad 2^j \cdot b$$

- If $i < j \Rightarrow 2^i \cdot b | 2^j \cdot b$.

- If $j < i \Rightarrow 2^j \cdot b | 2^i \cdot b$. \square

$$3. \quad X = \{1, 2, \dots, n\}$$

$$R = \{(x, y) : x \leq y\}$$

height = $n+1$

width = $\binom{n}{1}$

maximum = $\{1, 2, \dots, n\}$

minimum = \emptyset

56. Linear Orders

(Total/Linear Order): Let $P = (X, \leq)$ be a partially ordered set. We call P a total or linear order provided P does not contain incomparable elements.

(Isomorphic): We say $P \cong Q$

$$x \mapsto f(x)$$

"poset P is isomorphic to poset Q "

if:

isomorphism \rightarrow bijective invertible

1. there exists a bijection $f: P \rightarrow Q$

2. the bijection f is "order-preserving" such that

$$\forall x, y \in P, x \leq_P y \iff f(x) \leq_Q f(y)$$

1. Given $f: P \rightarrow Q$

isomorphism

Prove x is a minimum in P iff $f(x)$ is a minimum in Q .

Proof:

(\Rightarrow) Suppose x is minimum in P . (NTS: $f(x) \leq_Q z \forall z \in Q$)

1. By definition, $\forall y \in P, x \leq_P y$.

2. For any $z \in Q$, since f is surjective, $\exists y \in P$ such that $f(y) = z$.

3. By the isomorphism property of f , $\forall y \in P, f(x) \leq_Q f(y) = z$.

4. Therefore $\forall z \in Q, f(x) \leq_Q z$, making $f(x)$ minimum.

(\Leftarrow) Suppose $f(x)$ is minimum in Q . (NTS: $x \leq_P y \forall y \in P$)

1. By definition, $\forall z \in Q, f(x) \leq_Q z$

2. Let $z = f(y)$

3. Therefore $y = f^{-1}(z)$

4. Since f is an isomorphism, there exists an order preserving $f^{-1}: Q \rightarrow P$.

5. Applying f^{-1} to both sides of 1: $x \leq_P f^{-1}(z) = y$.

6. Since f^{-1} is surjective, this covers all the elements in P . \square

For what n do we get a linear order?

$n = p^k$ for some prime p
and nonnegative integer k .